## T-SGX: Eradicating Controlled-Channel Attacks Against Enclave Programs

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# The cloud is growing 7 times faster than the rest of IT

The latest IDC forecast says public cloud spending will grow almost 25% this year, topping \$122 billion. And the growth keeps up through 2020.

Network World | FEB 21, 2017 12:25 PM PT



#### 62 Percent of Companies Store Sensitive Customer Data in the Public Cloud

And almost 40 percent of cloud services are commissioned without the involvement of IT, a recent survey found.

By Jeff Goldman | Posted February 21, 2017



## **IT leaders say it's hard to keep the cloud safe**

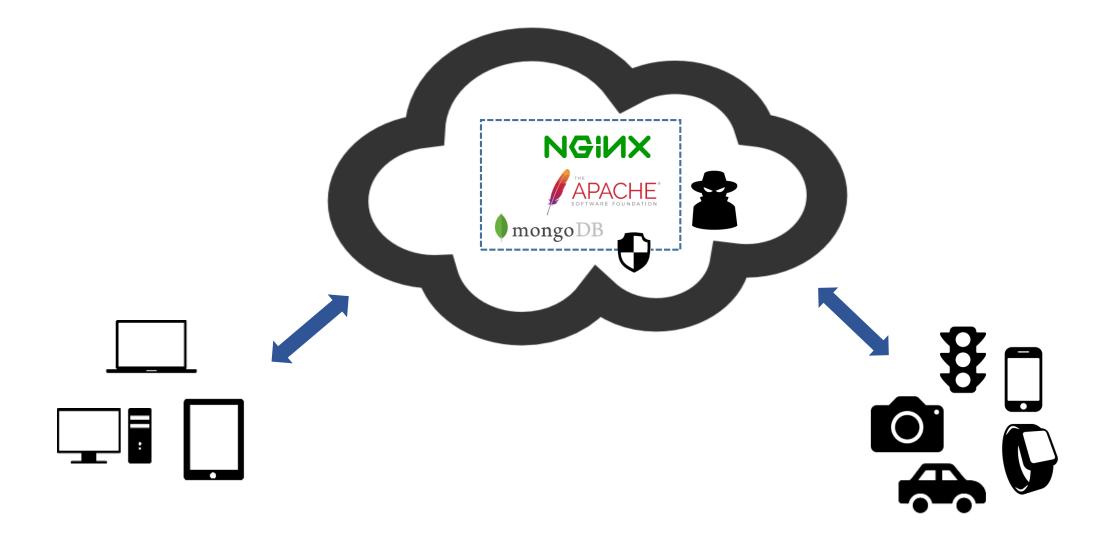
Shadow IT causing cloud trouble by illicitly working behind the scenes



**By Sharon Gaudin** | Follow Senior Writer, Computerworld | FEB 15, 2017 12:17 PM PT

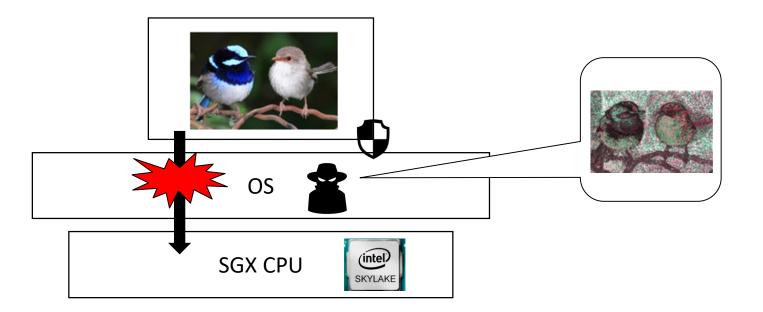


#### Intel SGX aims to secure users' code and data in the cloud



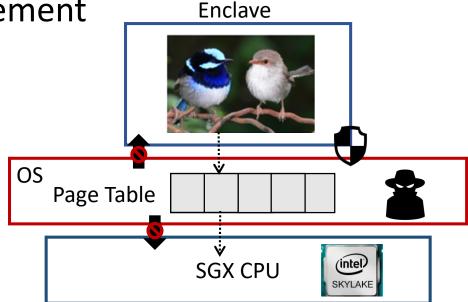
#### Controlled-channel attack [Oakland 2015] raises concerns

- An accurate side-channel attack that extracts the SGX-protected data
- Compromise the security guarantees of SGX



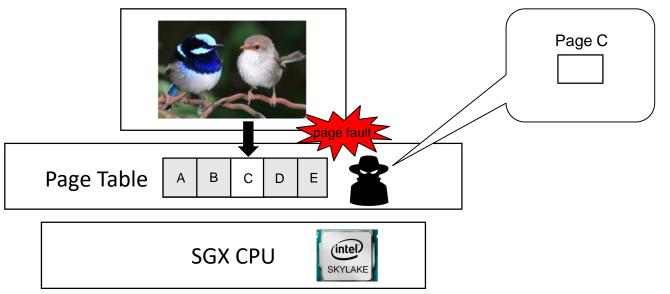
#### How the attack works (1/3)

- Intel SGX protects *enclaves* against an untrusted OS
- SGX still relies on the OS for resource management (e.g., memory mapping)



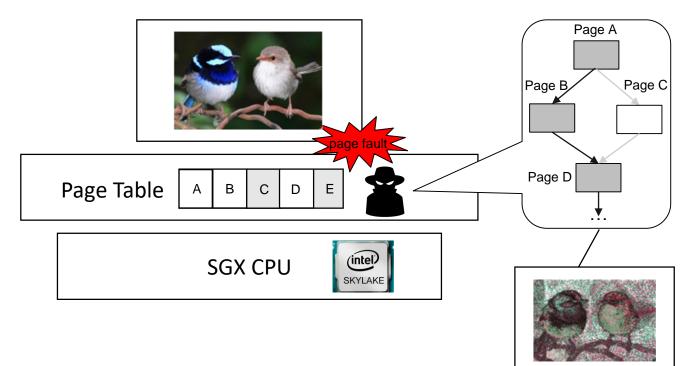
## How the attack works (2/3)

- Attacker fully controls the OS
- Page-fault side channel
  - Step 1: Unmap a page
  - Step 2: Enclave accesses the page
  - Step 3: Observe a page fault



## How the attack works (3/3)

- If the program's memory accesses depend on a *secret*, then this *secret* is being leaked
- Attack steps
  - Offline analysis
  - Obtain page-fault sequence
  - Infer the secret

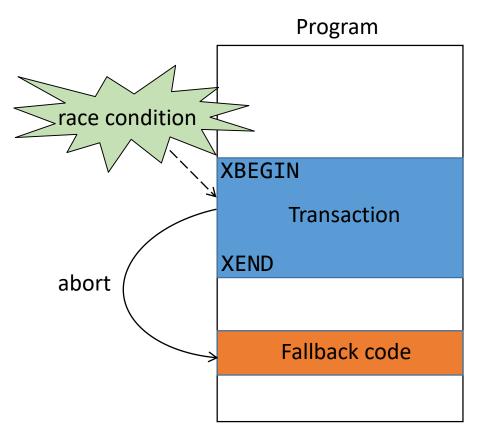


#### T-SGX Goals

- Prevent the controlled-channel attack
- The design should be practical
  - No hardware modification
  - Reasonable performance
  - Minimal developer effort (no need for program rewritten)

#### Intel TSX

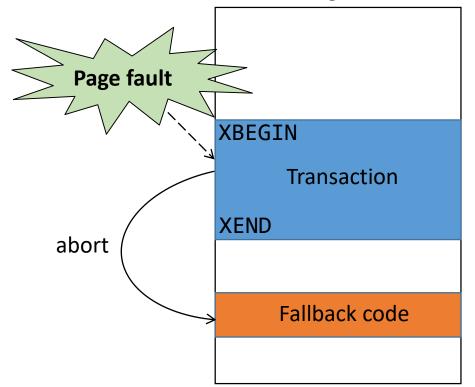
- CPU extension present in all recent Intel CPUs (since 2013)
- Supports hardware transactional memory
- Race conditions cause transaction abort
- An abort triggers fallback execution
  - Rolls back all changes
  - Control transfers to the fallback point



### Idea: Intel TSX also suppresses page faults

CPU does not deliver page faults to the OS
Page faults to th

• OS *cannot observe* the page fault inside a transaction



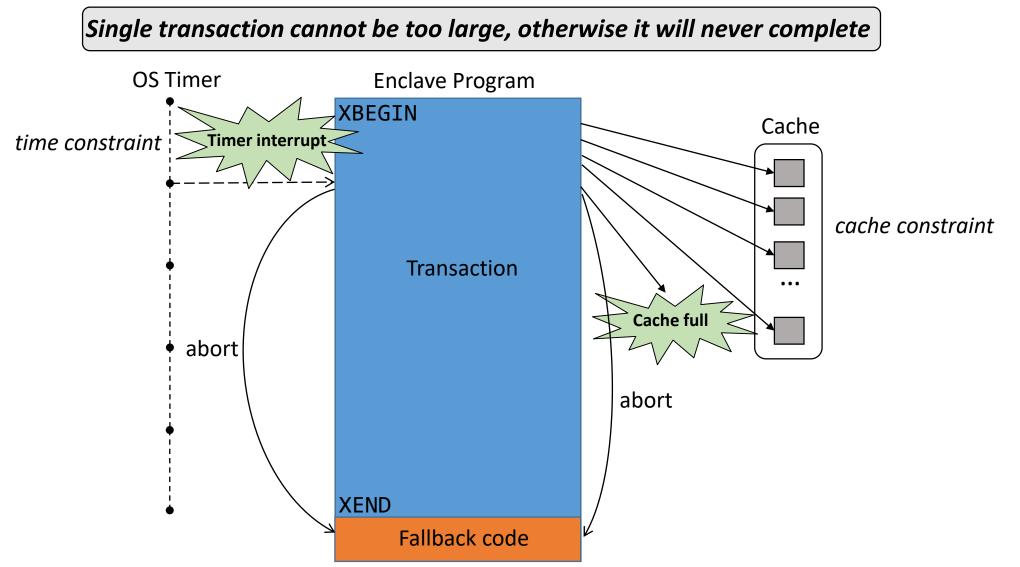
Program

#### The strawman design

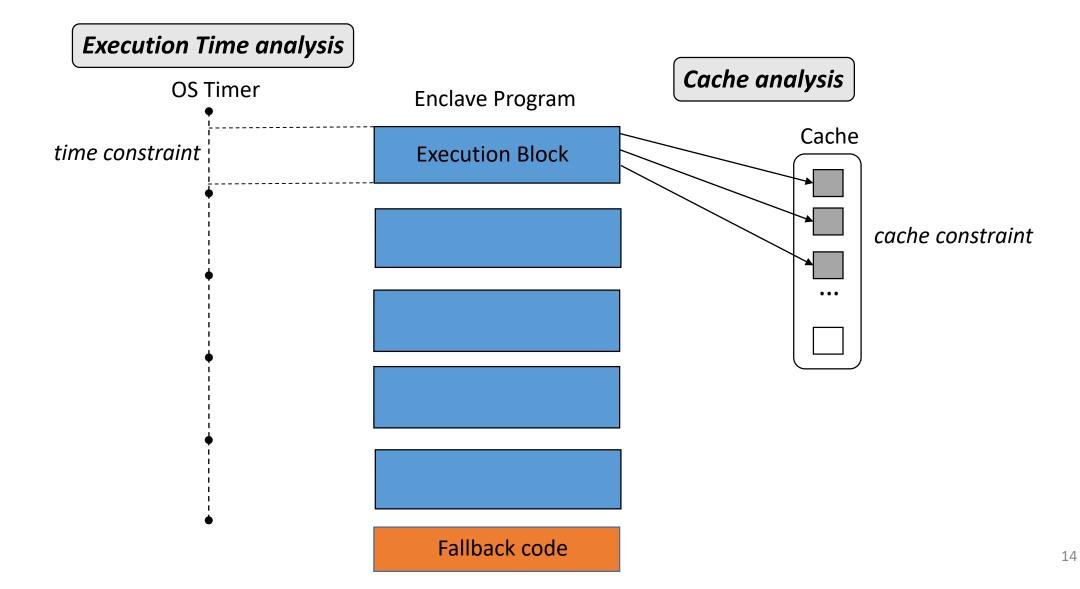
 Make the whole enclave **Enclave Program** Page fault as a transaction XBEGIN • Enable the self-detection to page faults inside the enclave Transaction abort XEND

Fallback code

## Challenges



#### Solution: Break a program into execution blocks



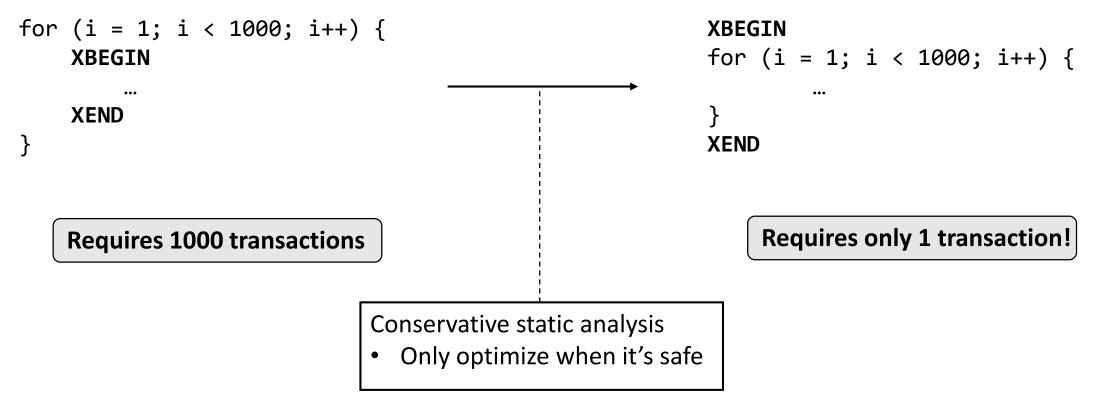
## Optimization: merging tiny blocks (1/2)

• **Problem**: Setting up transaction comes with a fixed cost (~200 cycles)

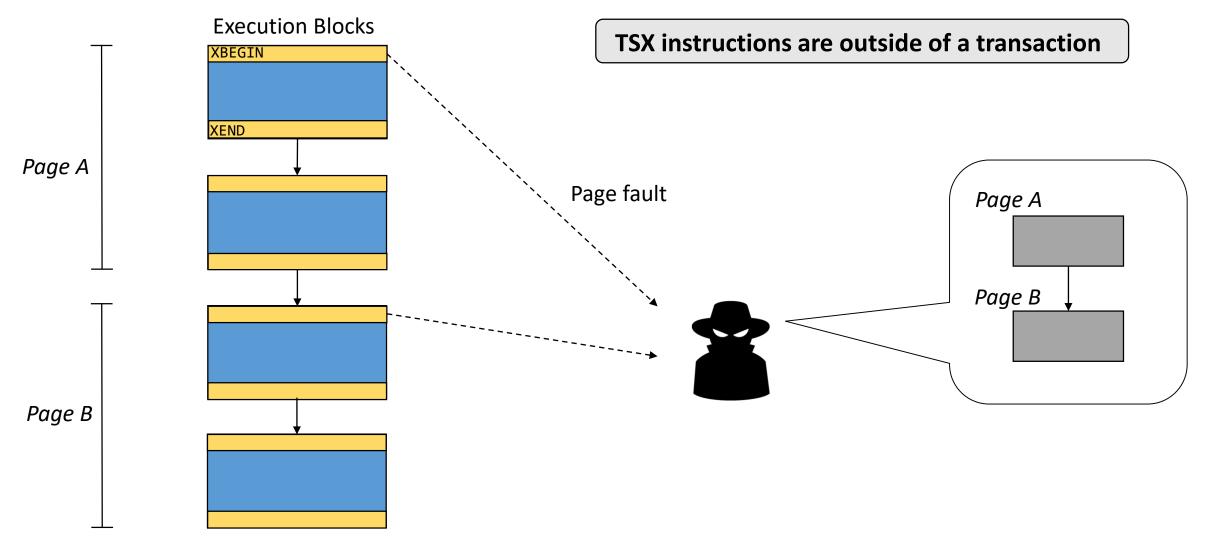
- If continuous blocks satisfy the cache and time constraints, we merge them
  - Loops
  - If-else statement
  - Functions

### Optimization: merging tiny blocks (2/2)

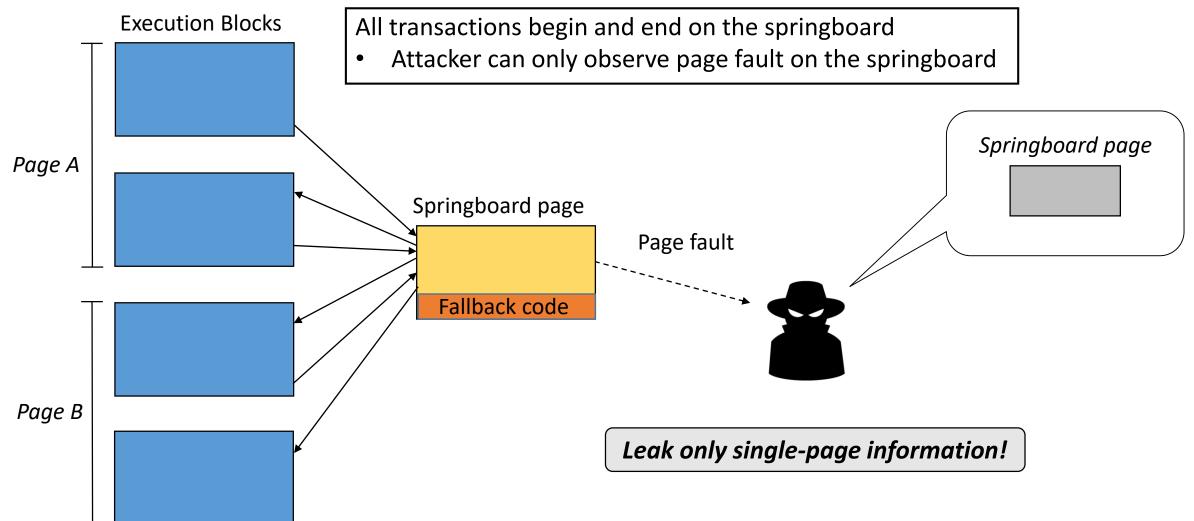
• Example: Loop optimization



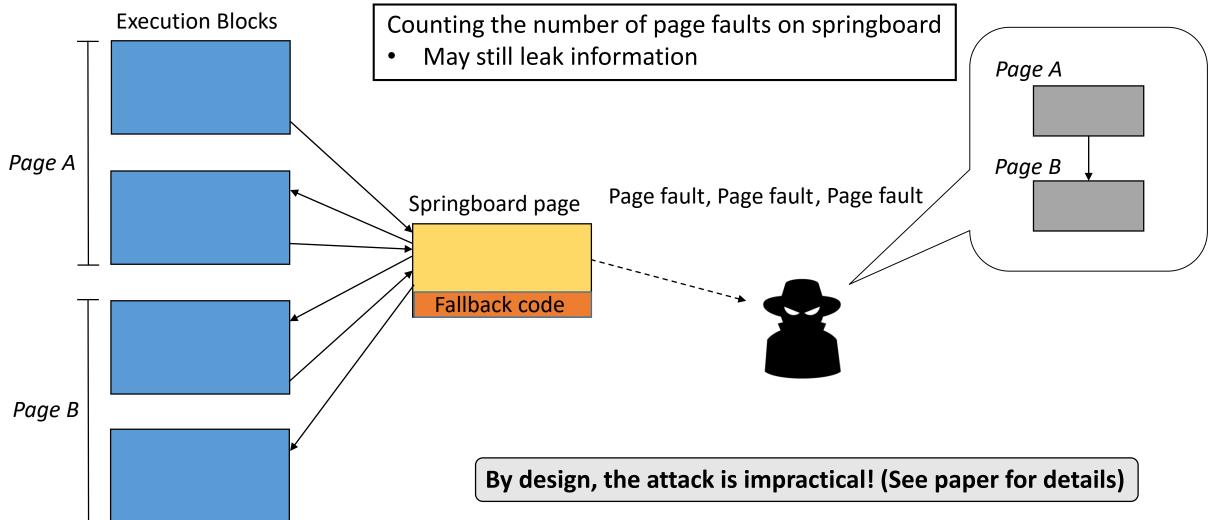
### This design still leaks information



### Solution: Springboard design



#### Springboard design also prevents advanced attacks



#### Implementation: T-SGX

- Based on the LLVM compiler
  - Mostly modifying LLVM backend
  - 4,100 line of code
  - Fully automated program transformation

#### Evaluation

- How general is the T-SGX approach?
- How much overhead does a transformed program have?

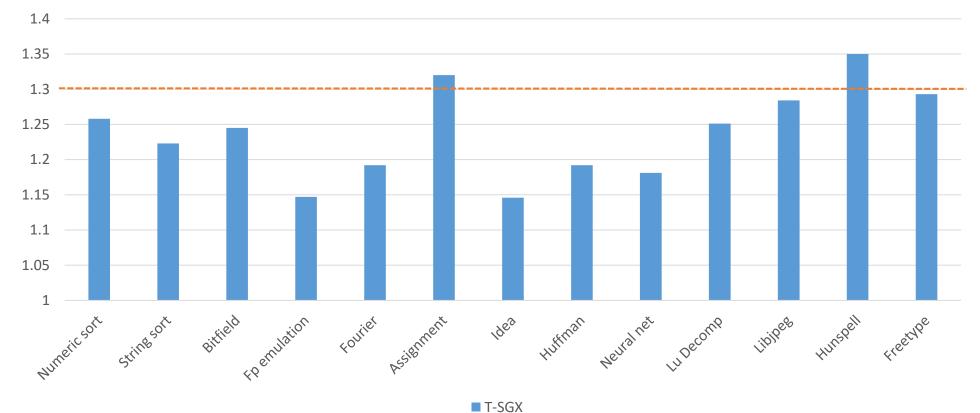
## T-SGX works for general C/C++ programs

- 0 lines of source code change
- Fully-automated compiling chain

Application	Line of Code
Numeric sort	211
String sort	521
Bitfield	225
Fp emulation	1,396
Fourier	235
Assignment	490
Idea	353
Huffman	448
Neural net	746
Lu decomposition	441
Libjpeg	34,763
Hunspell	24,794
FreeType	135,528

#### T-SGX incurs reasonable overhead

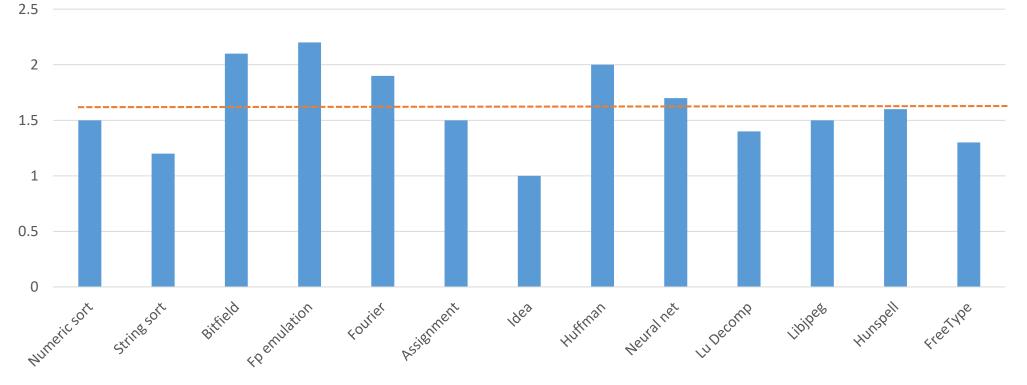
- Average 30% memory overhead
  - Additional instructions for each execution block



Benchmark programs

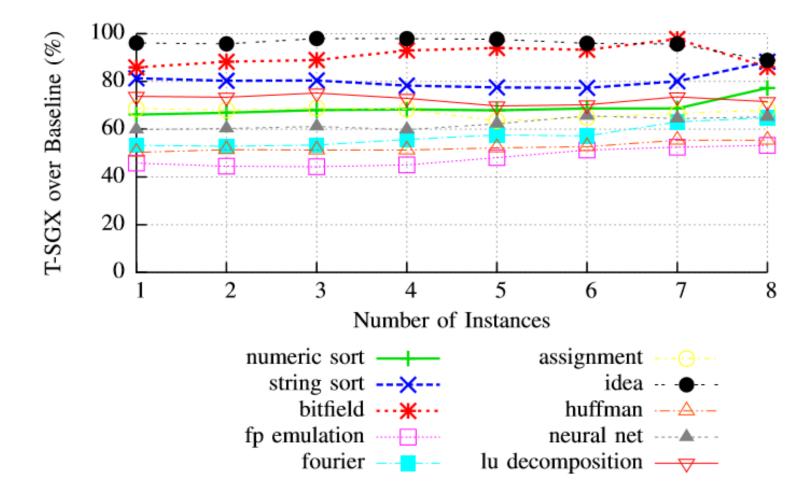
#### T-SGX incurs reasonable overhead

- Average 50% runtime overhead (geometric mean)
  - Largely depends on number of loop iterations that repeatedly start a transaction



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#### Consistent runtime overhead on concurrent execution



#### Conclusion

• We proposed and implemented T-SGX, which effectively protects enclaves against the controlled-channel attack.

#### • T-SGX

- Requires no hardware modification
- Incurs reasonable runtime overhead and still has potential to improve (e.g., using more advanced program analysis or performance profiling)
- Automatically transforms a program without the need for manual effort

#### Q&A