Asynchronous intrusion recovery for interconnected web services

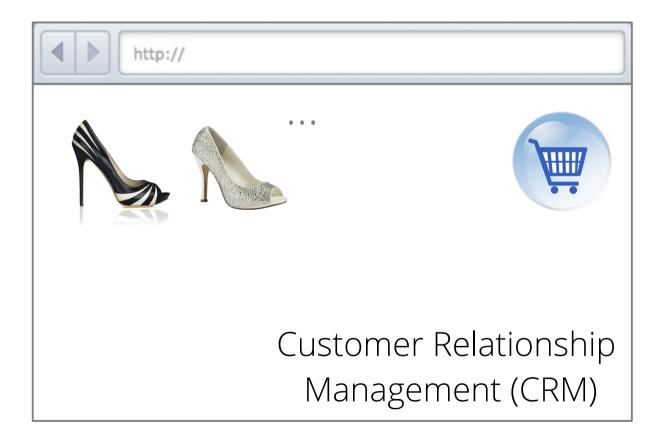
Ramesh Chandra, Taesoo Kim, Nickolai Zeldovich

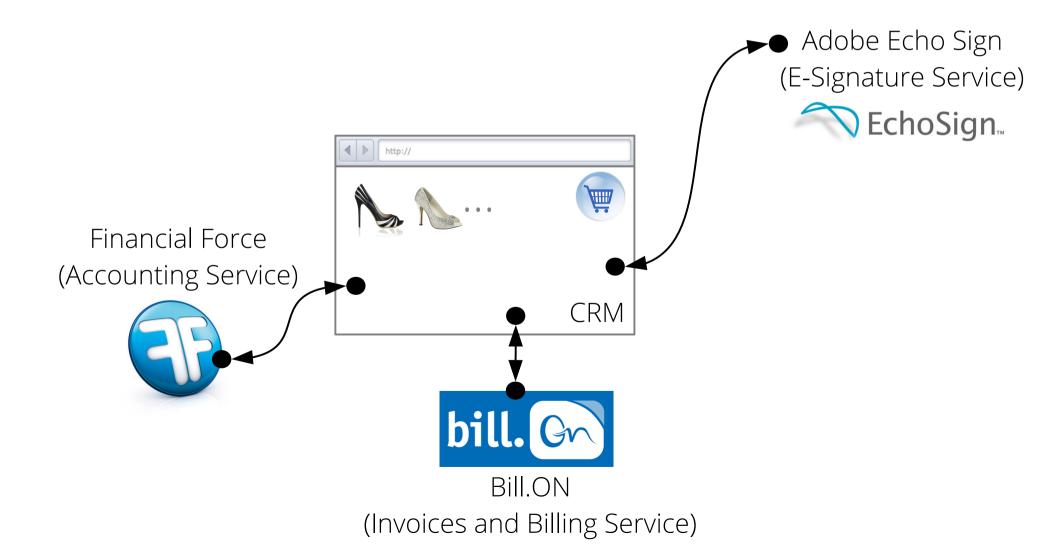
MIT CSAIL

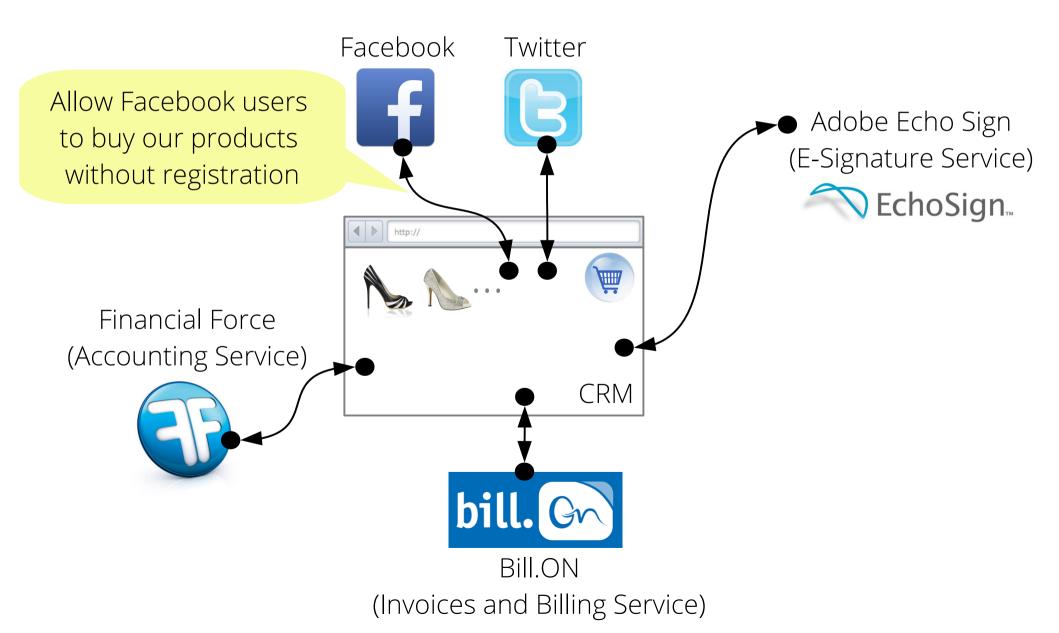
Today's web services are highly interconnected

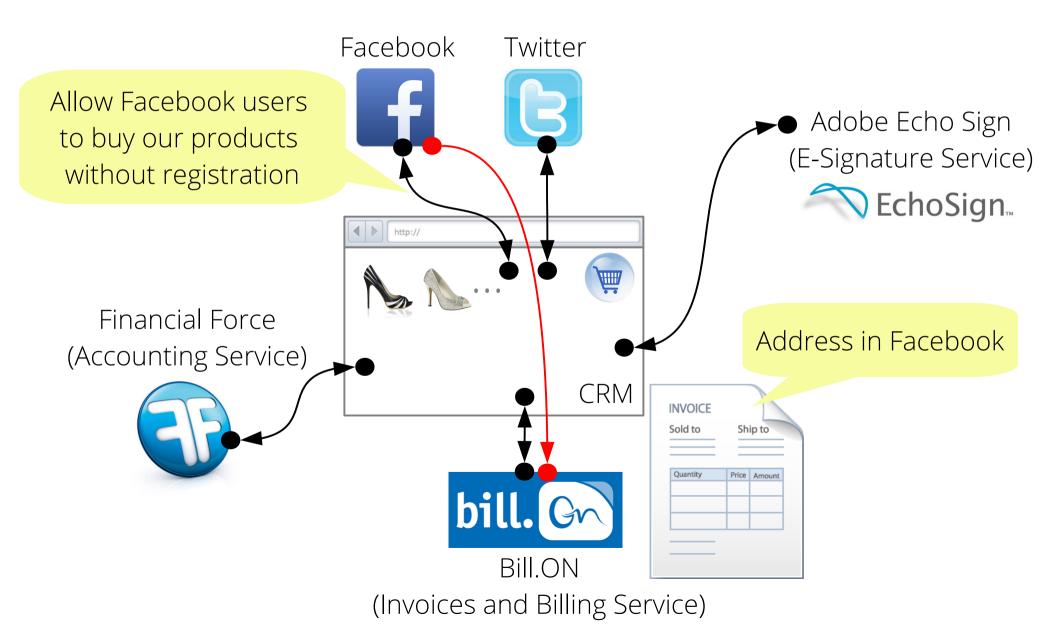
• Many web services provide APIs to other sites

- Many websites integrate those APIs:
 - Authentication: Facebook Connect, Google+ ...
 - Data sharing: Dropbox ...
 - Business process management: Salesforce ...

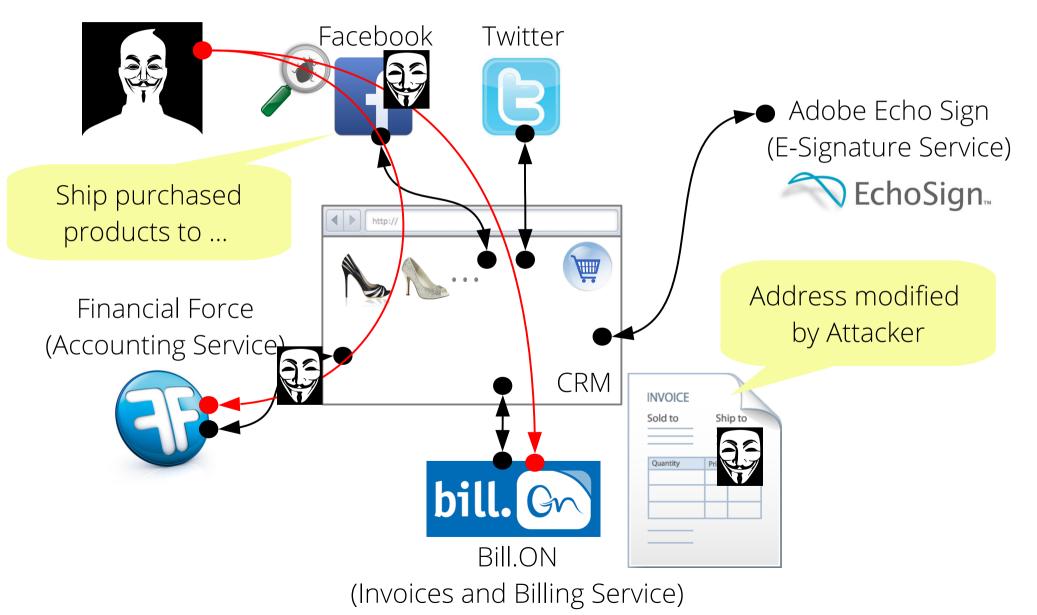








Attack in one service can spread between services



Bugs in web services are commonplace

- Facebook (Mar 29th 2013):
 - Attackers can intercept full permission access tokens



Bugs in web services are commonplace

- Facebook (Mar 29th 2013):
 - Attackers can intercept full permission access tokens
- Many web services have similar bugs
 - Twitter (Aug 20th 2013)
 - Instagram (May 2nd 2013)
 - Microsoft Yammer (Aug 4th 2013)

SECURITY



Goal

- **Recovering integrity** in interconnected services
 - Repair the state of affected services as if the attack never occurred

- State-of-the-art: manual recovery
 - Admin doesn't trust other sites for recovery
 - Require manual interaction (e.g., email other admin)

General plan for automatic recovery

- Use **rollback-and-replay** for recovering integrity in single machine
 - Prior works: Retro [OSDI '10], Warp [SOSP '11]

• Extend rollback-and-replay to many web services!

Challenges

- Rollback-and-replay requires **global coordinator**
 - Each service cannot decide what to do for repair

- All services must be **available** during recovery
 - We want to repair some services even if others are down
 - Consistency problem: some services are not repaired yet

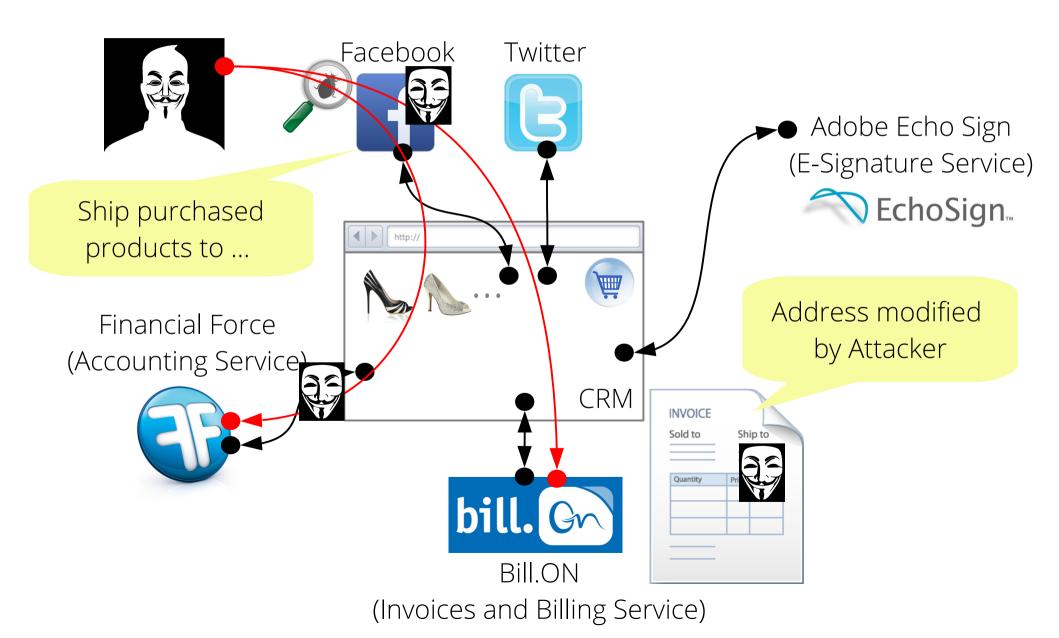
Contributions

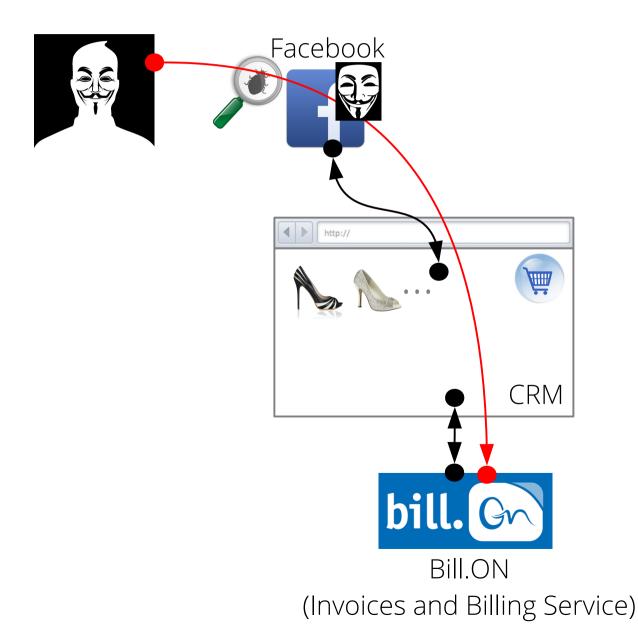
Enable **automatic** intrusion recovery in **distributed** web services

- 1. Repair protocol between services
 - No central coordinator
 - Each service controls its repair

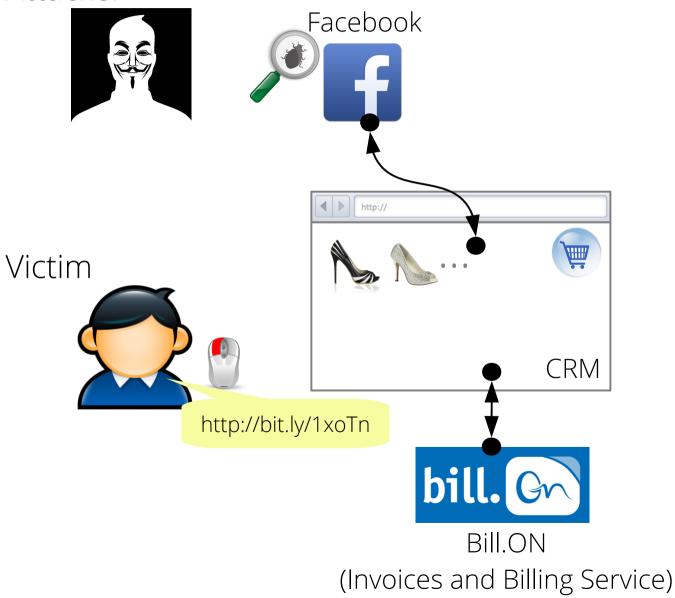
2. Asynchronous repair

- Proceed repair even with unavailable services
- Consistency in partially repair state

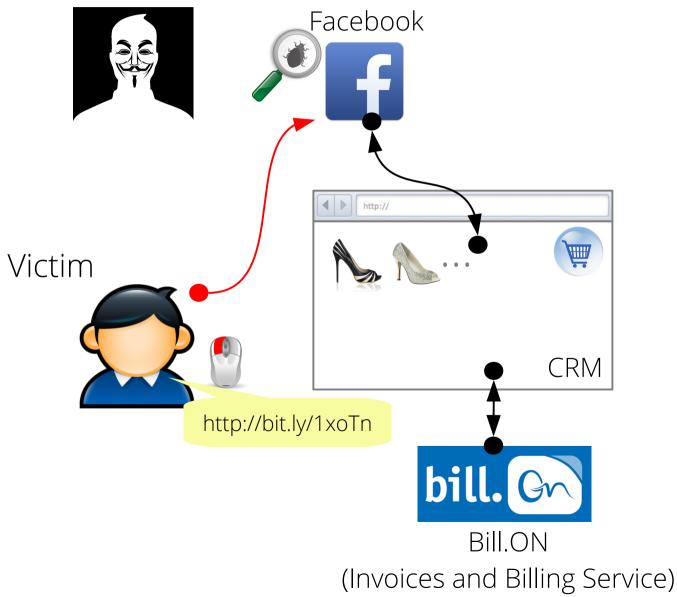




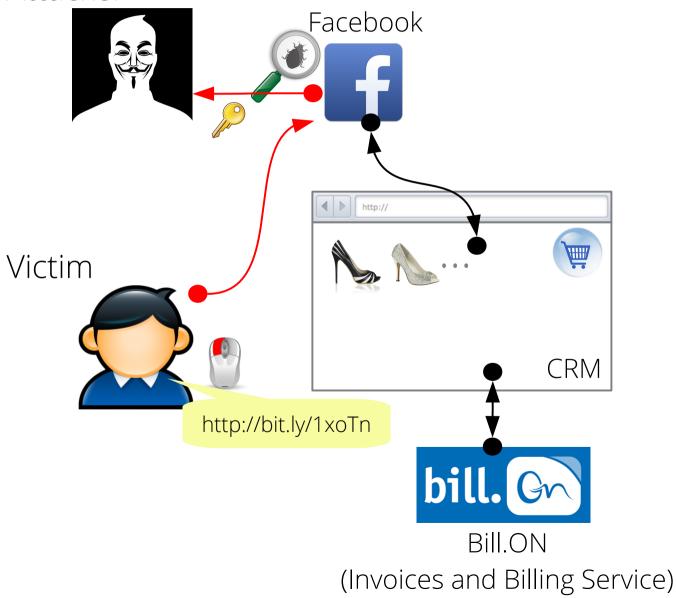
Attacker

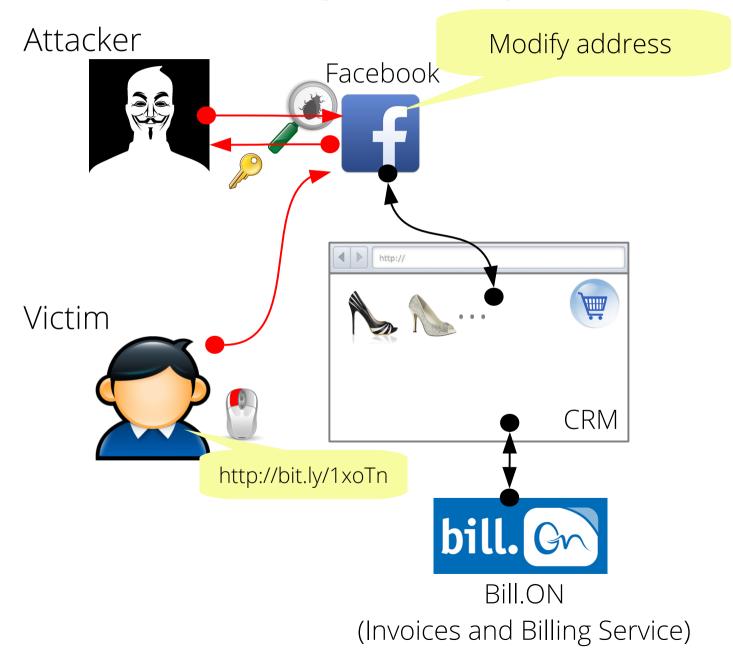


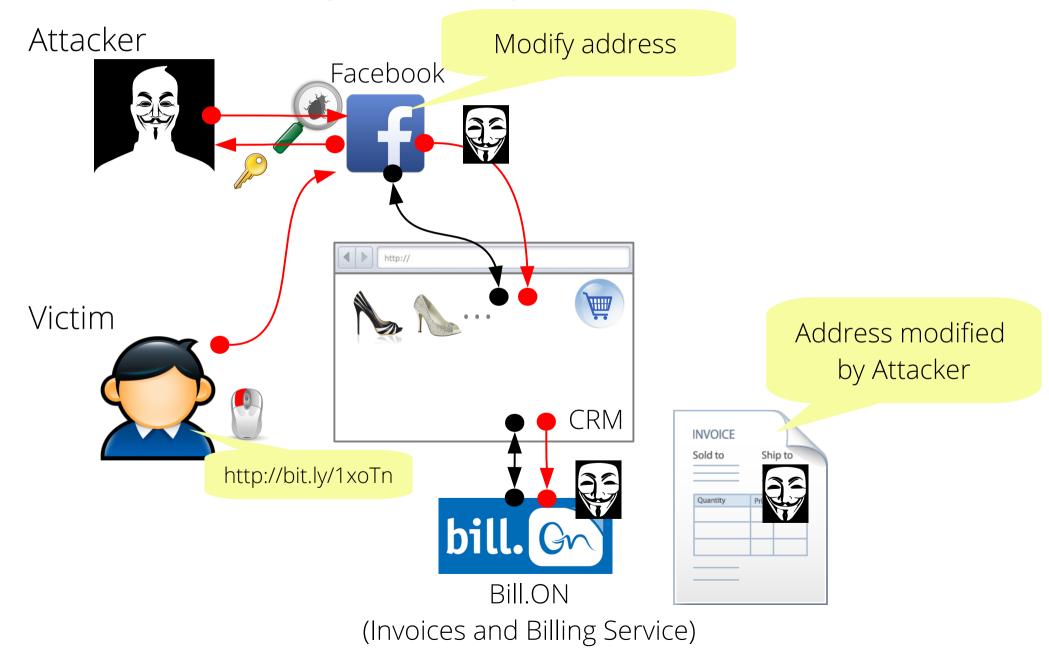
Attacker



Attacker





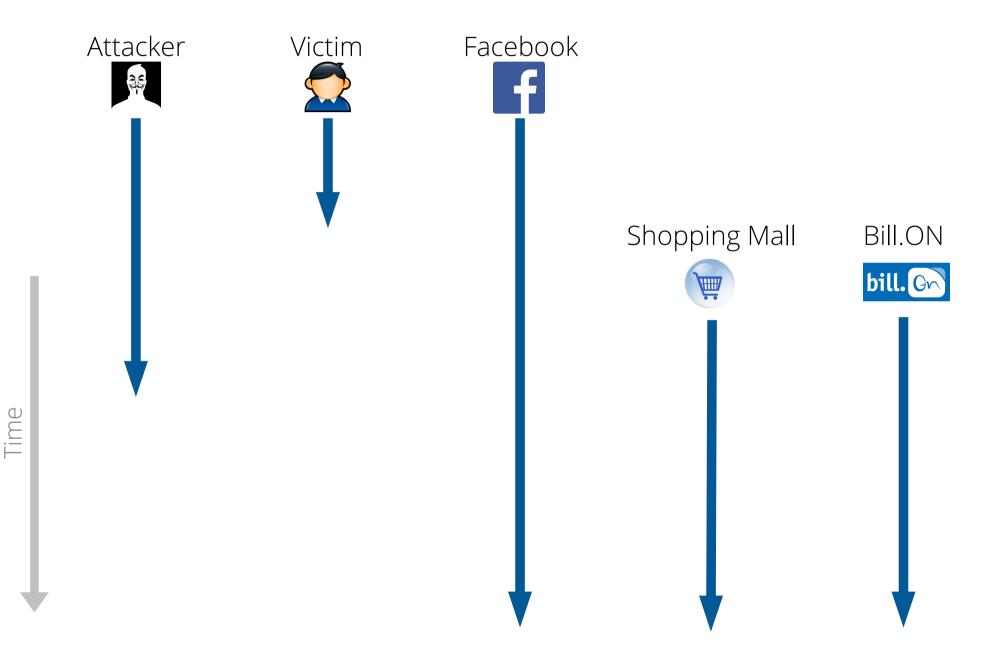


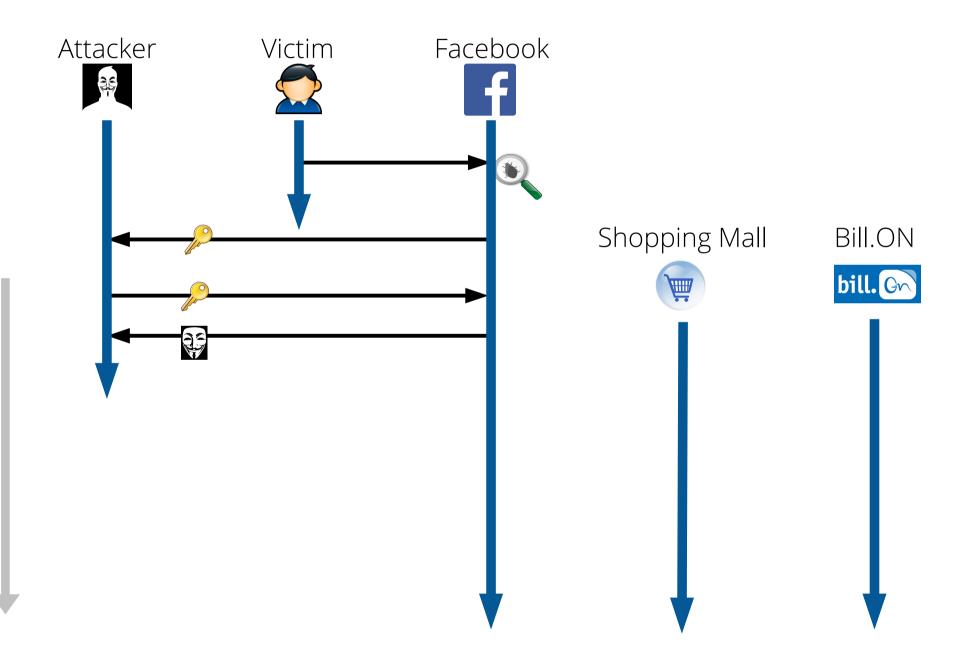


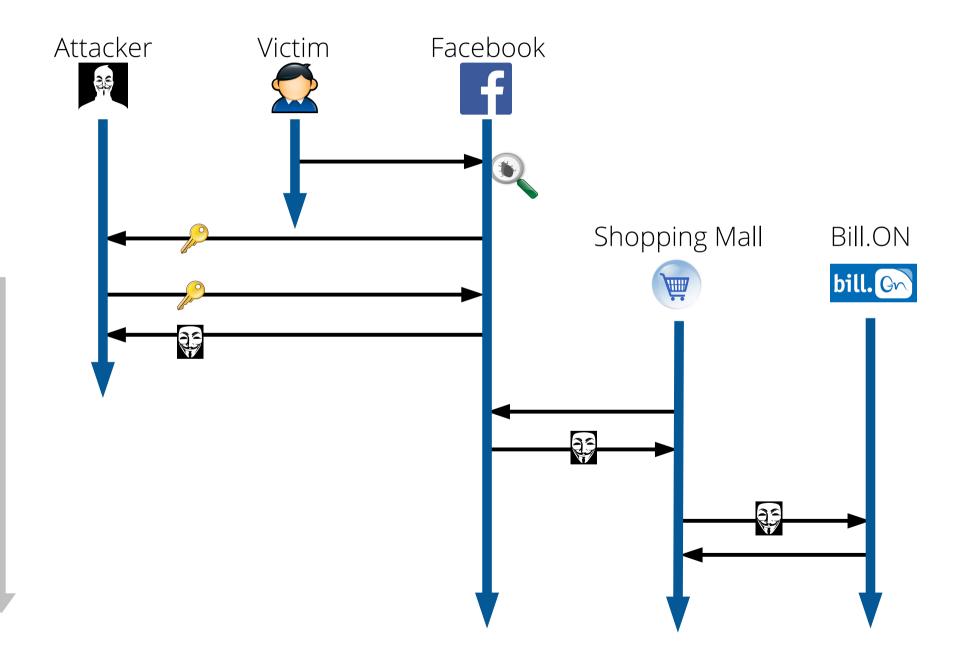




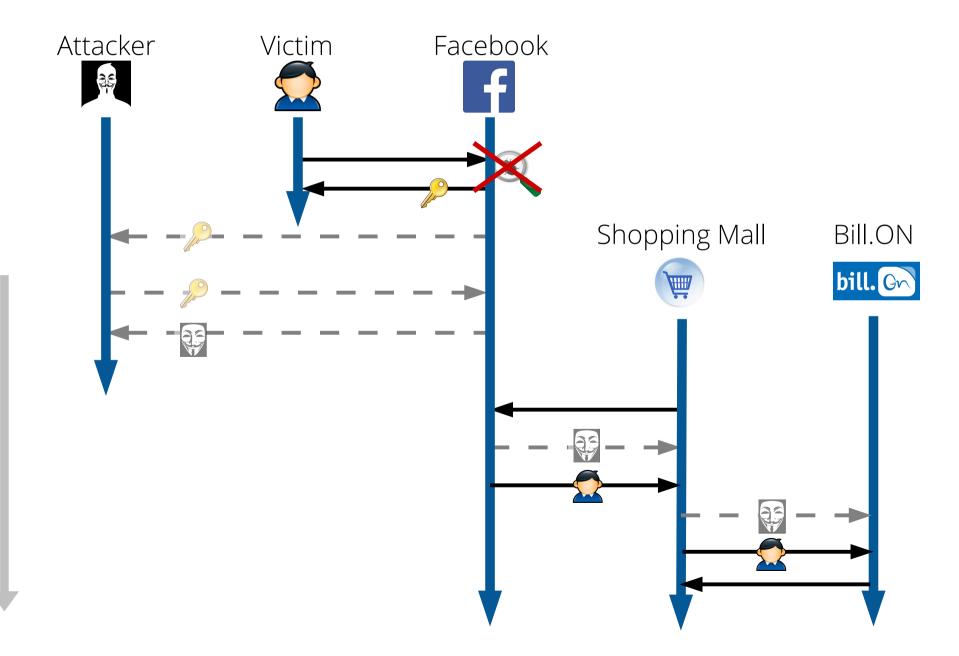




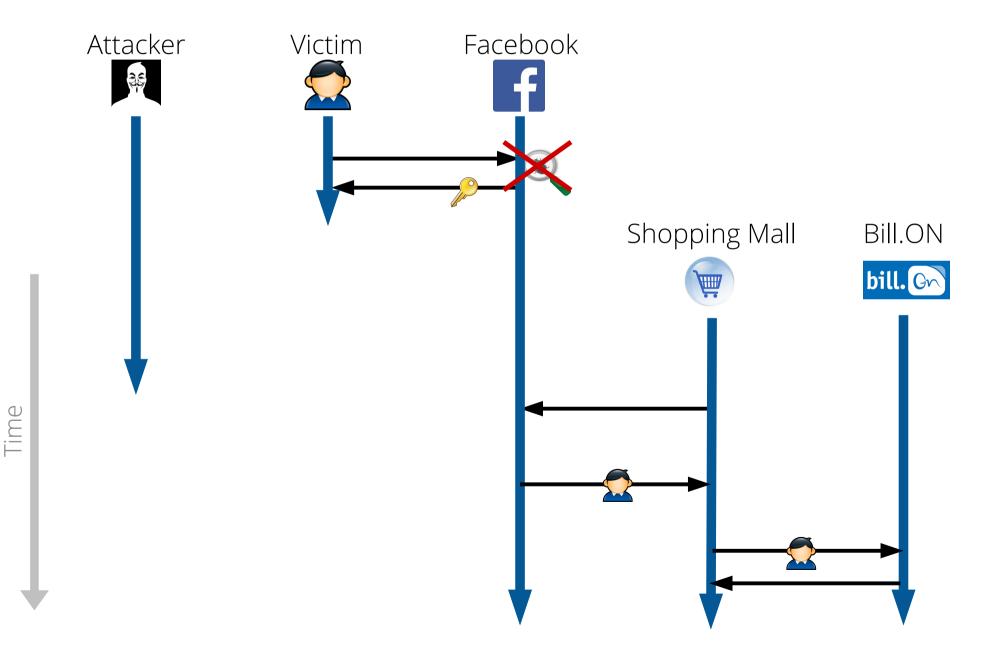




Goal: attack did not take place



Goal: attack did not take place



Overview of system execution

- Normal execution:
 - Record enough information for rollback-and-replay

• Repair:

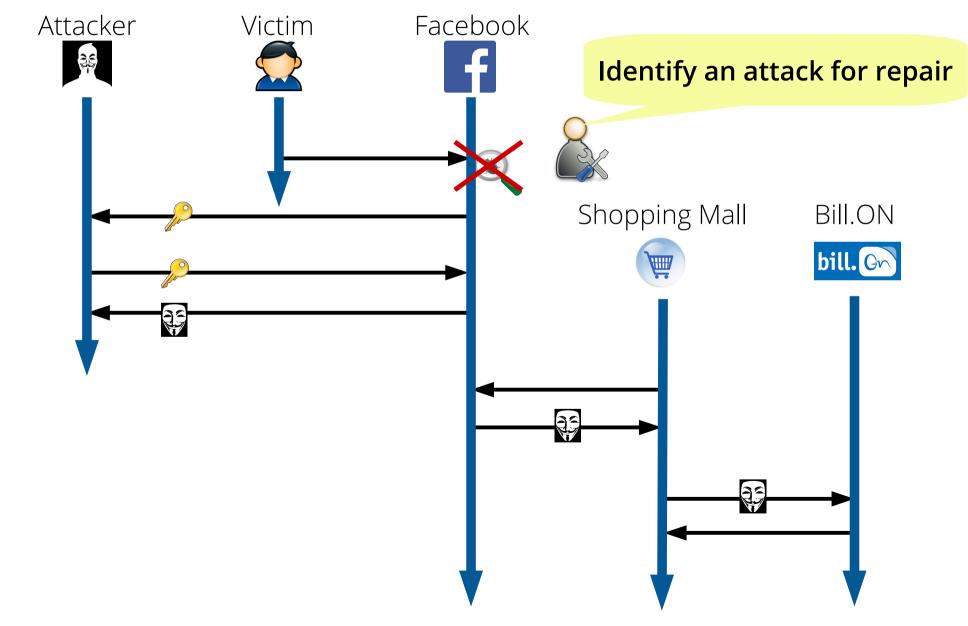
- Identify an attack to initiate repair
- Repair local state: rollback and replay recorded requests
- Propagate repair whenever local repair affects others

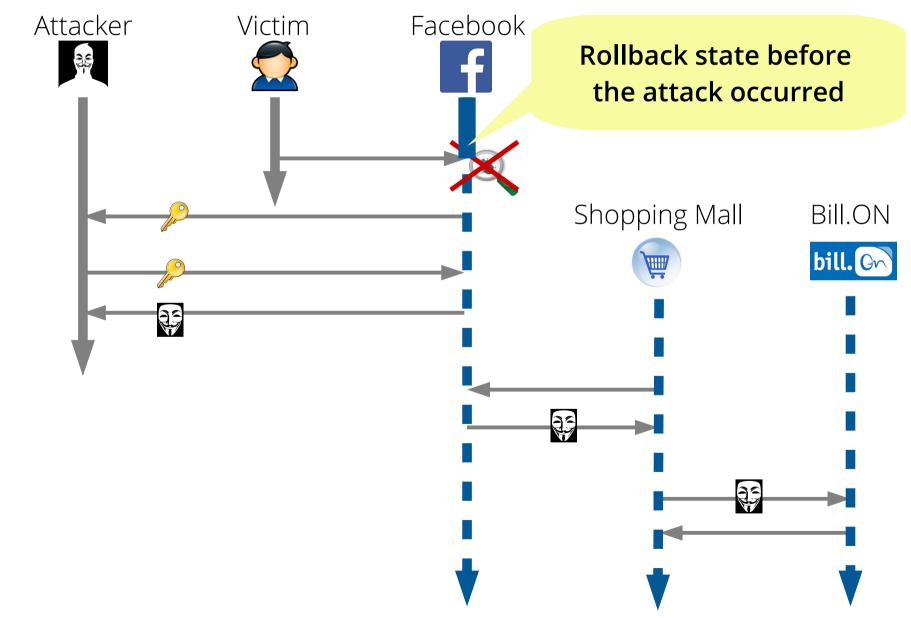
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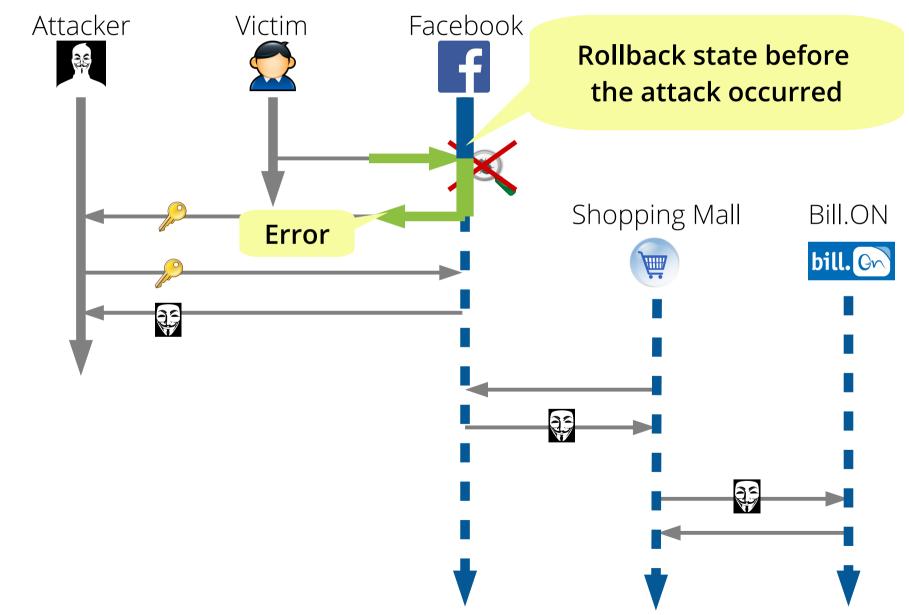
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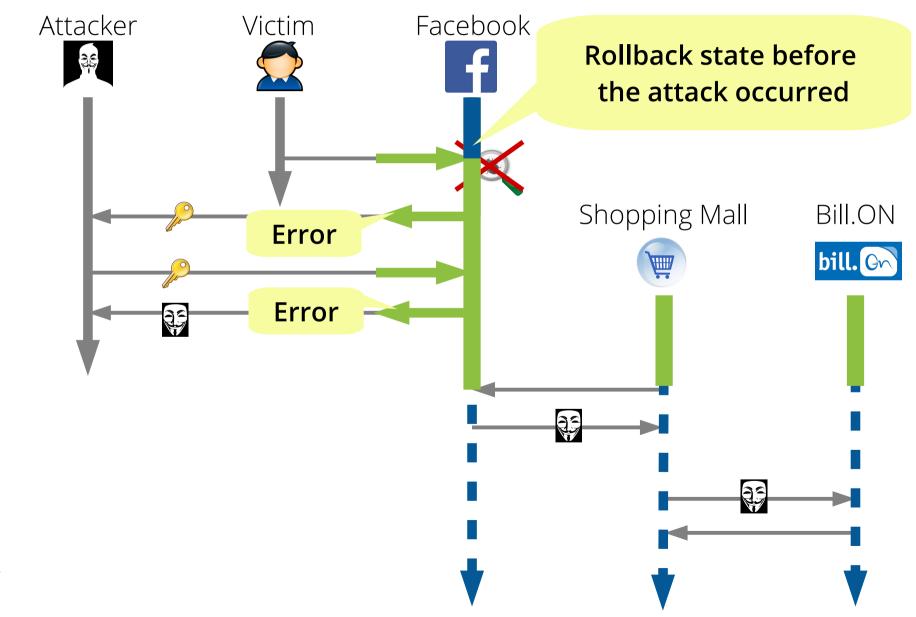
• Repair:

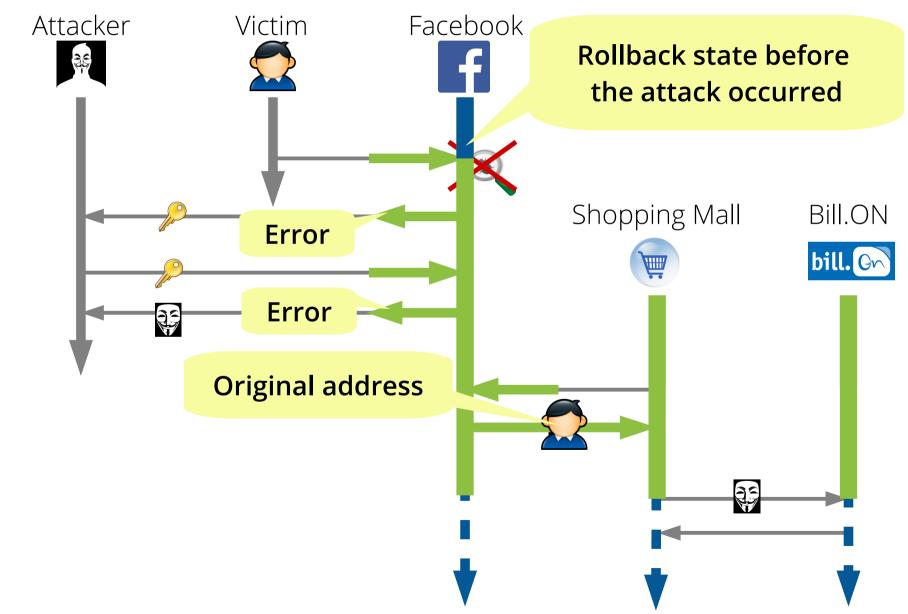
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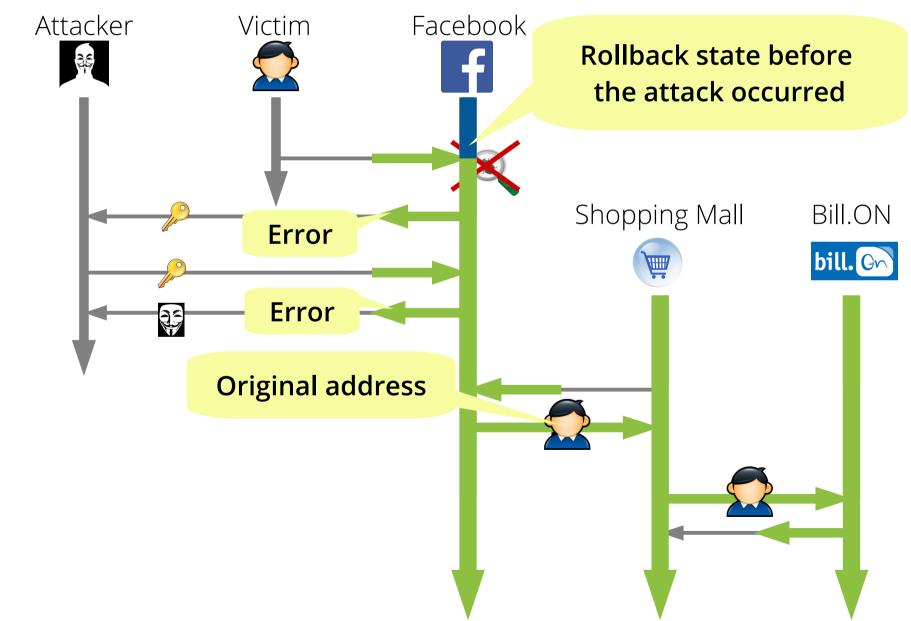


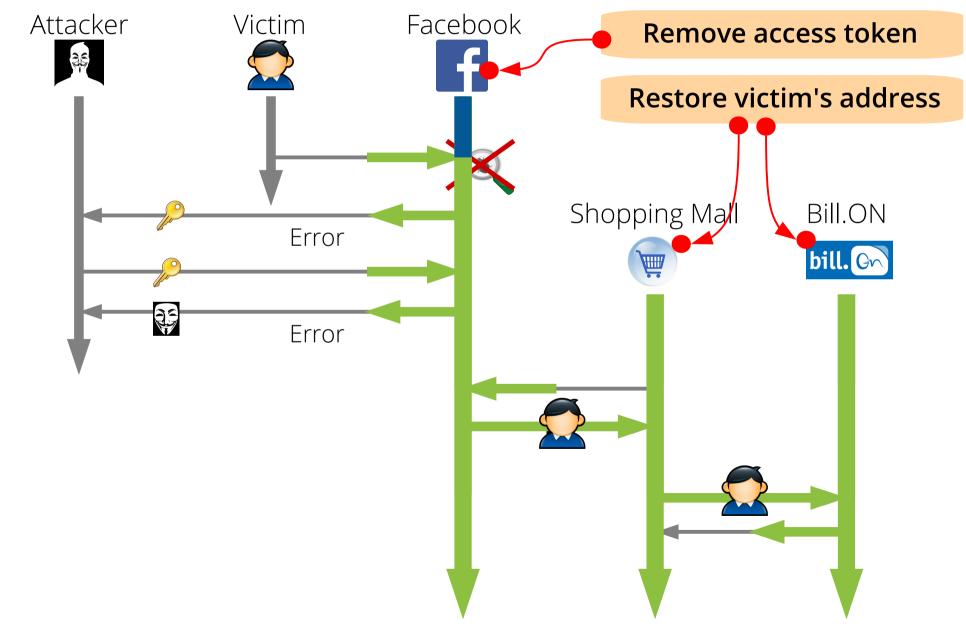












Problems in Strawman design

• P1. All services must be available

→ Support asynchronous repair with speculation

• P2. Require global coordinator

→ Define repair APIs between services

Problems in Strawman design

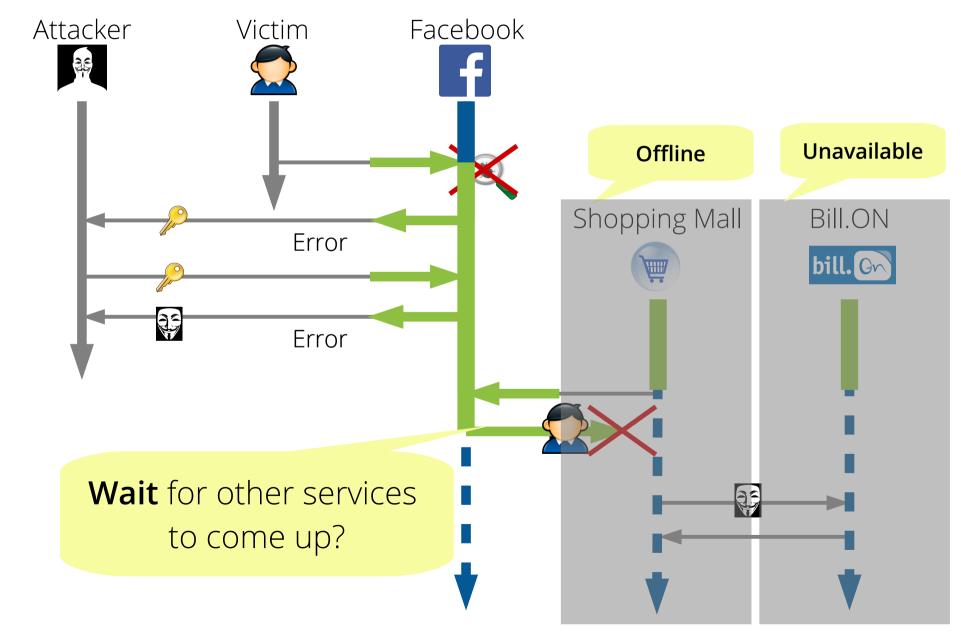
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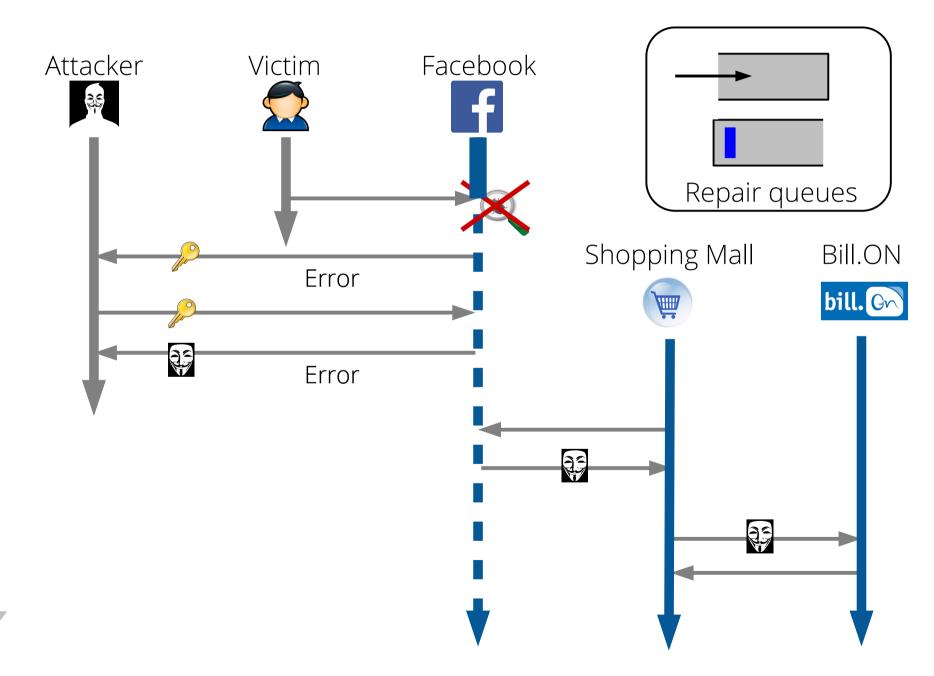
Challenge: cooperating with unavailable web services



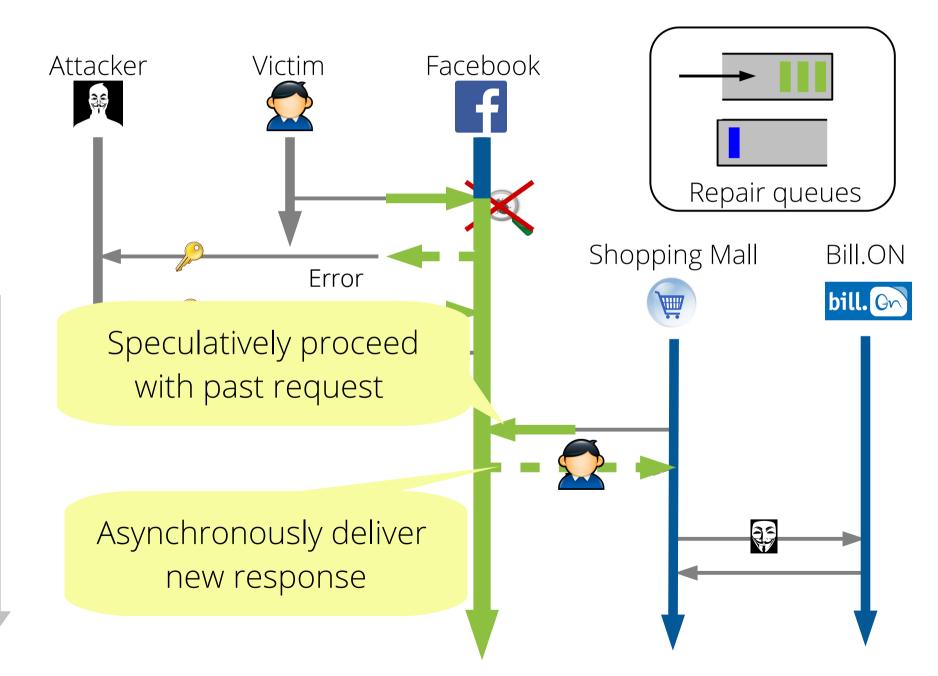
Solution: asynchronous repair

- Asynchronously deliver repair requests
- Speculatively proceed local repair with past responses (or timeout responses)
- Expose repaired state after local repair

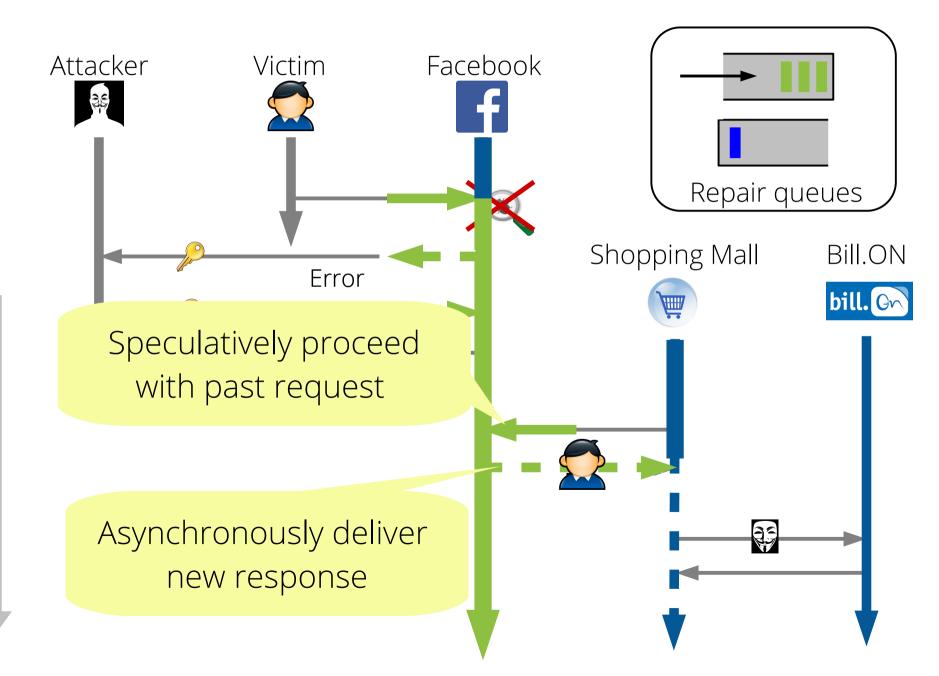
- Intuition: why asynchronous repair works?
 - Many web services are designed for independent operation, prepared for handling others failures

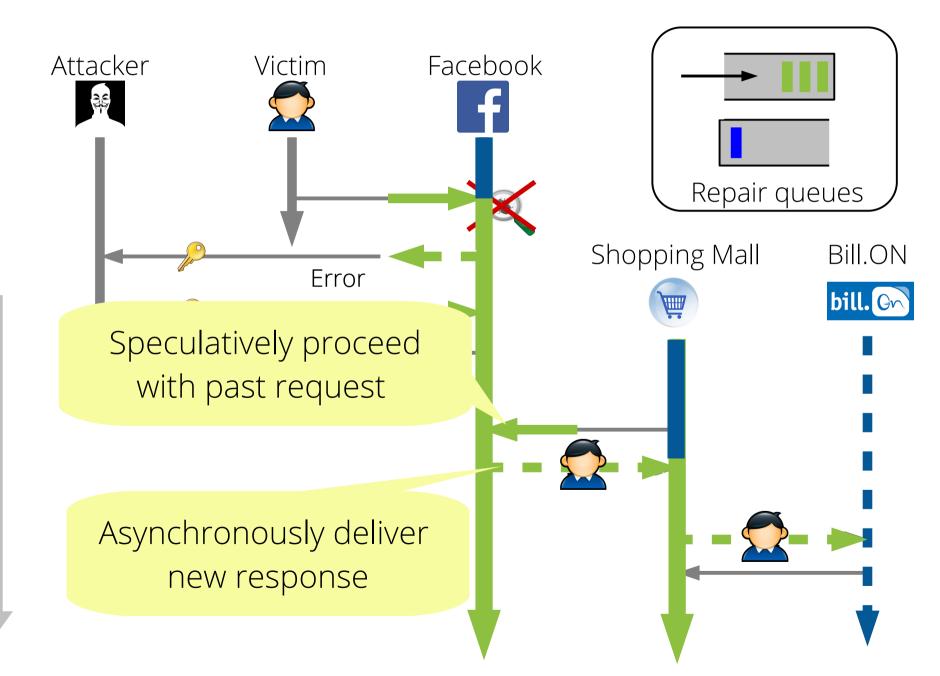


Time



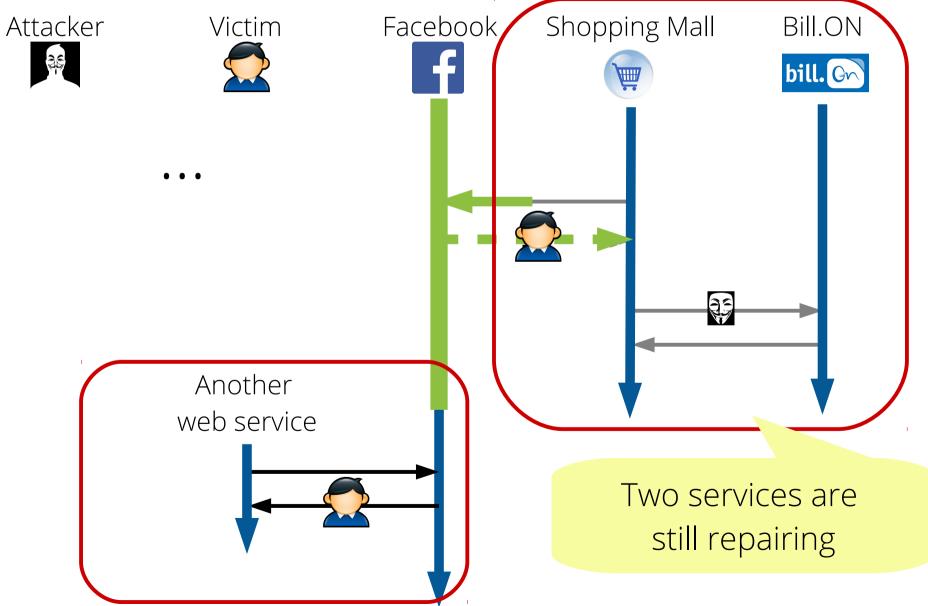
Time





Time

Example: exposing state after local repair

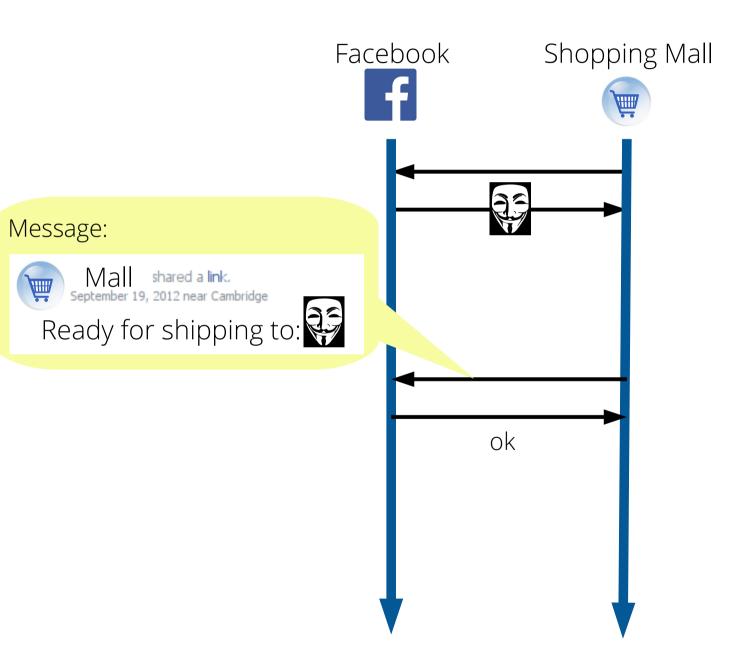


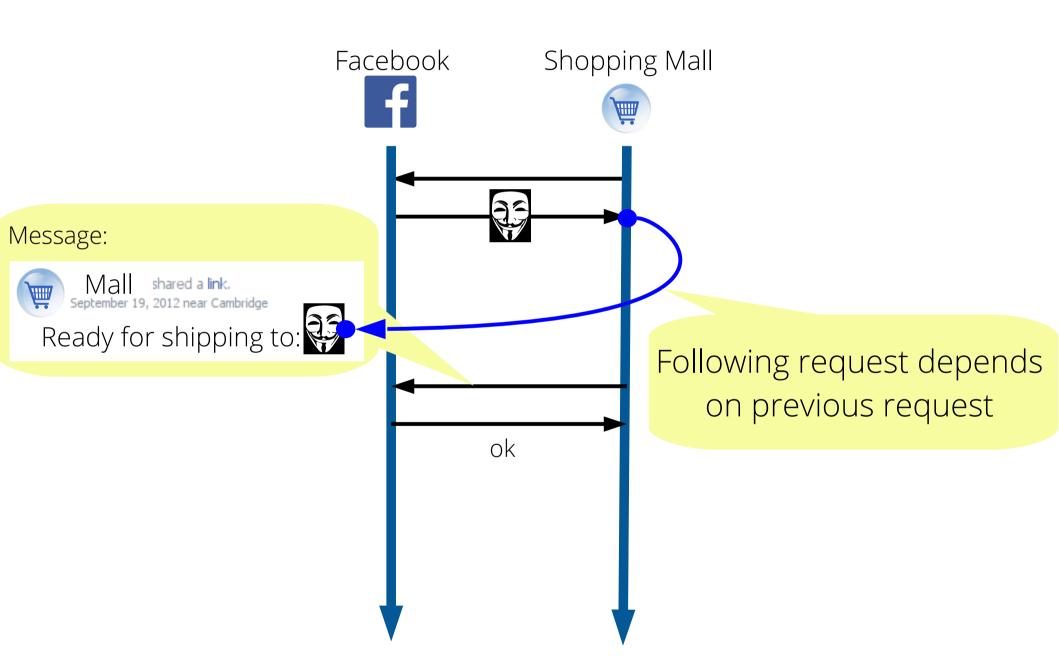
Time

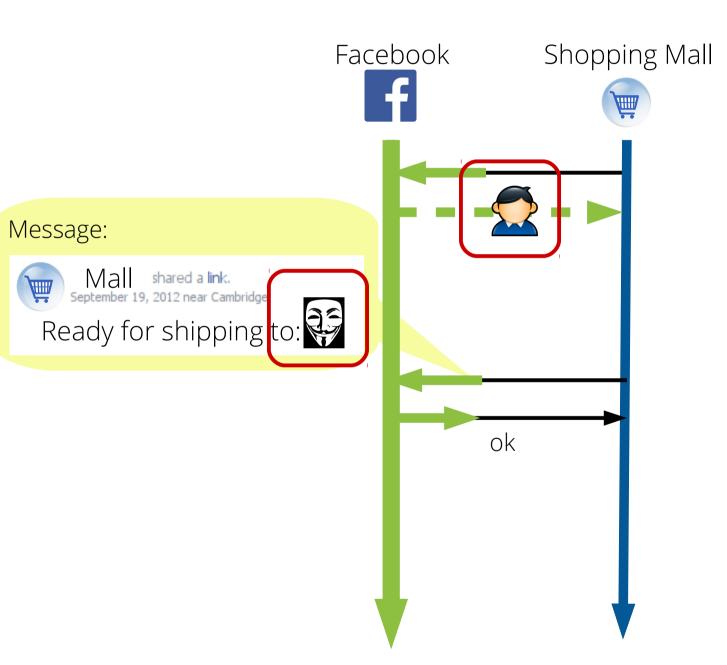
What if speculation fails?

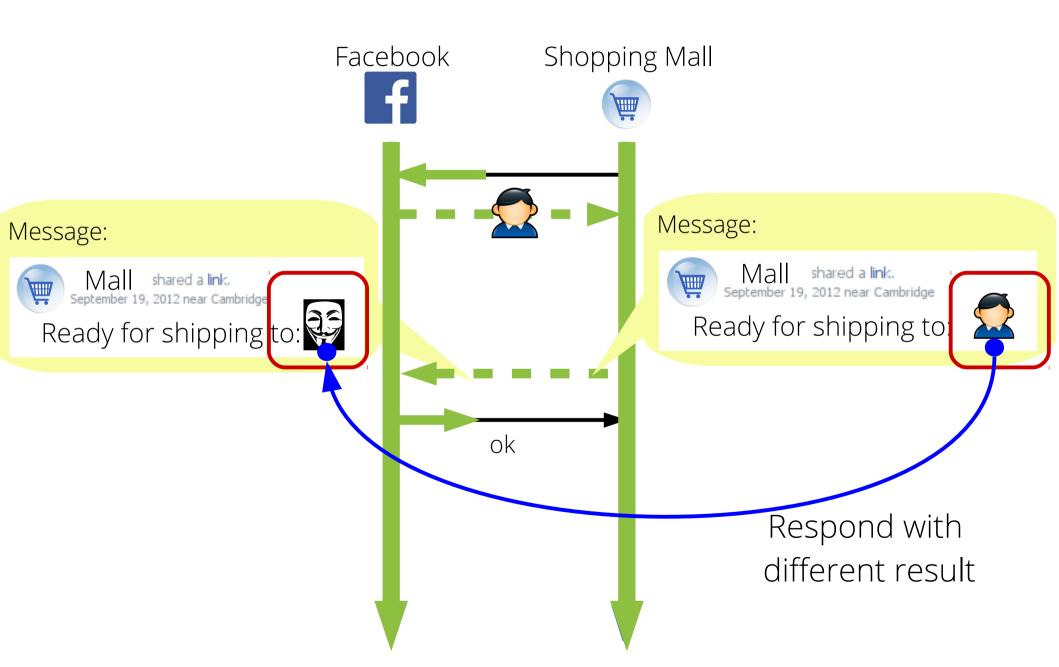
- If service responds differently,
 - Restart local repair with the new response
 - In fact, it is not different from initiating new repair

• Asynchronous repair will converge to the correctly repaired state at the end

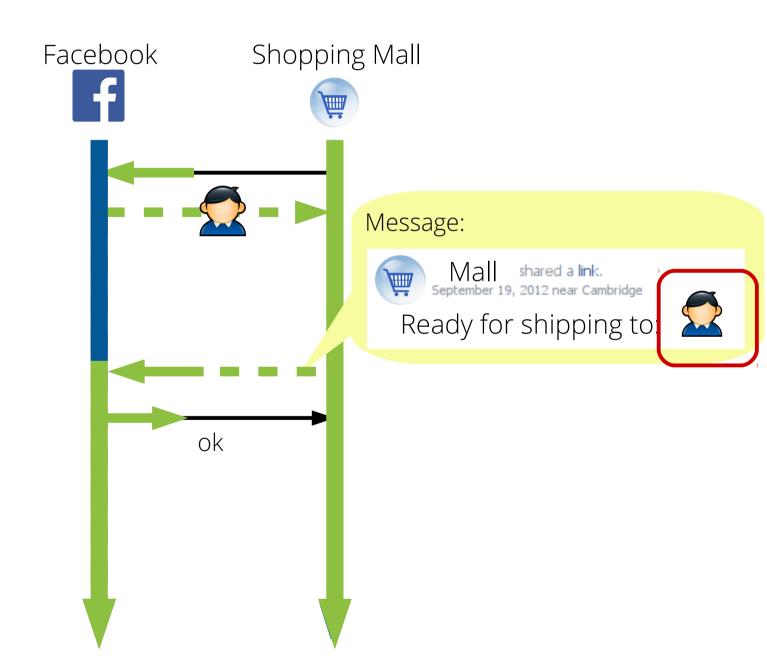








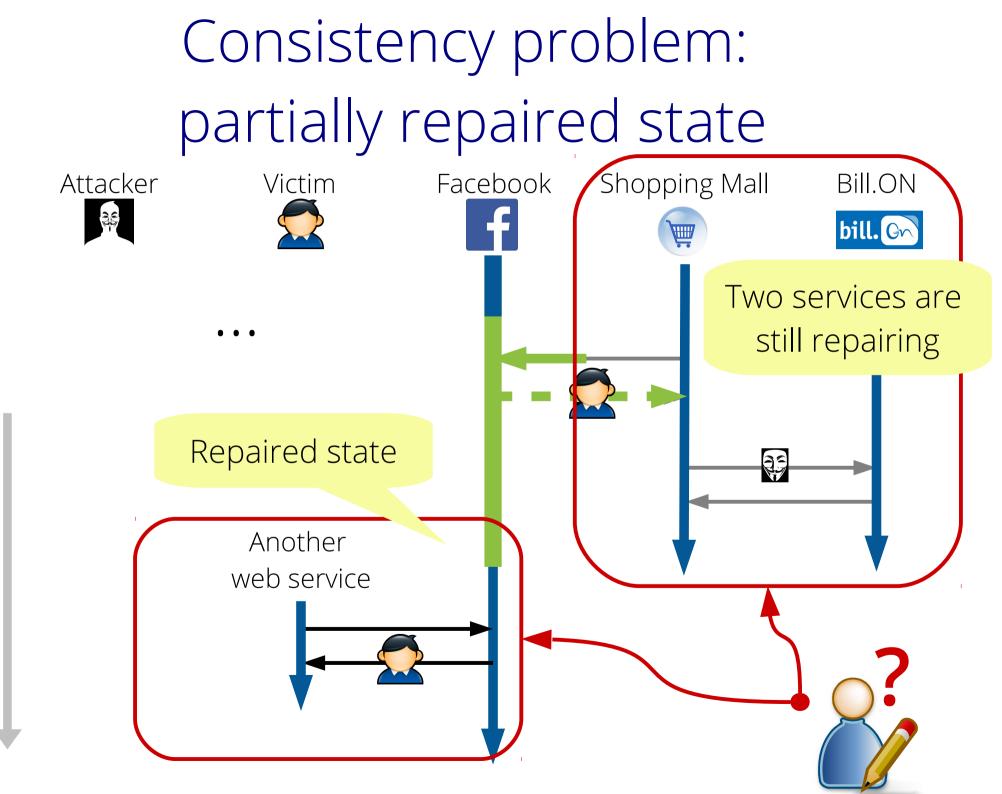






Asynchronous repair makes **forward progress** in time-line graph, so it will **converge** to the correctly repaired state at the end

ok



Time

Consistency problem: partially repaired state

- Exposing partially repaired state might diverge global state
 - But it is not something new that our recovery mechanism introduces more
 - Most of web services already cope with this problem

Exposing partially repaired state might violate service invariants

- **Service invariants**: guarantees by service provider (e.g., locking serivce: when lock is held, no concurrent access)
- In theory: yes (for arbitrary tightly coupled systems)
- In practice: no
 - RESTful APIs usually provide consistency per API
 - Web services are in nature **loosely coupled**

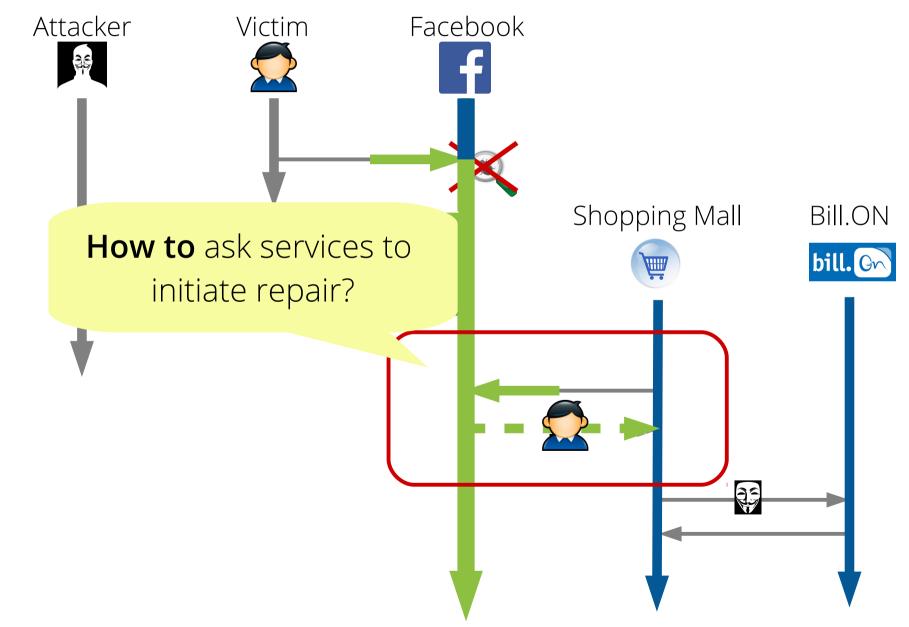
Consistency: partial repair state

- Services and clients **already deal with concurrency**
- Repair of a service is modeled as:
 - Being performed by a concurrent **repair actor client**
 - That uses the service's regular API calls
- So, partially repaired state can be considered as state resulting from yet another concurrent operations

Problems in Strawman design

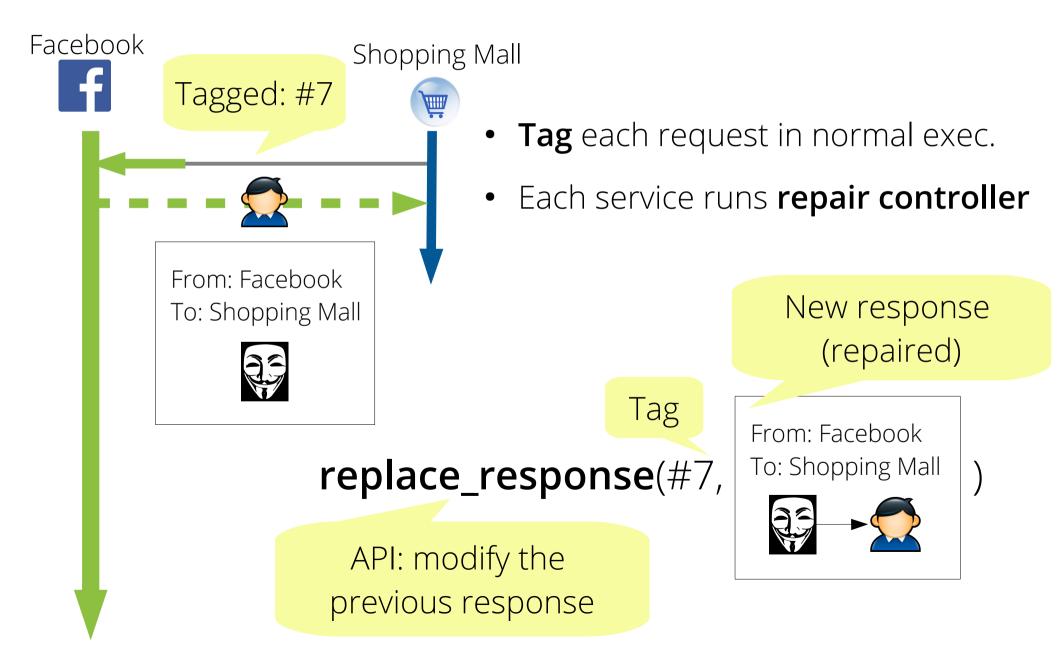
- P1. All services must be available
 - → Support asynchronous repair with speculation
- P2. Require global coordinator
 → Define repair APIs between services

How to propagate repair requests without global coordinator?



Time

Requesting repair with APIs (RPC)



Repair APIs (RPC)

- No centralized coordinator, each server invokes following repair APIs to recover from the attack
 - **replace_response**(tag, resp): replace past response
 - replace_request(tag, req): replace past request
 - **delete**(tag): delete past request
 - **create**(req, before, after): execute new requests in the past

Repair APIs (RPC)

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 - **delete**(tag): delete past request
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If service supports those 4 APIs, it can participate in decentralized recovery

Authentication of repair APIs

- Too application specific
 - (e.g. Email service: sender can delete recipient's emails?)
- **Delegate** authentication to original web services
 - Implement application specific policy (e.g. ask admin for confirmation of repair)
 - Assign a credential to repair requests

Summary of design

1. Asynchronous repair

- Proceed repair with offline or unavailable services
- Consistency in partially repair state

2. **Repair APIs** between services

- No central coordinator
- Each service controls its repair
- Delegate authentication

Implementation

- Prototype implementation: **Aire**
 - Extend **Django** web framework
 - Support existing Django app. with **few modifications**
 - Support Askbot, Django-OAuth, and Dpaste
 - e.g., Askbot's authentication policy: 55 LoC
 - Total: 5700 lines of Python code

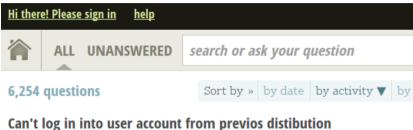
Evaluation questions

- Can Aire support real web services?
- Can Aire recover from distributed attacks?
- What are the runtime overheads of Aire?

Aire supports real web services Askbot OAuth Provider

Dpaste

Askbot



#fedora19#login user old

Multiboot UEFI WIN8/Fedora19: GRUB not starting ==> Fedora not booting

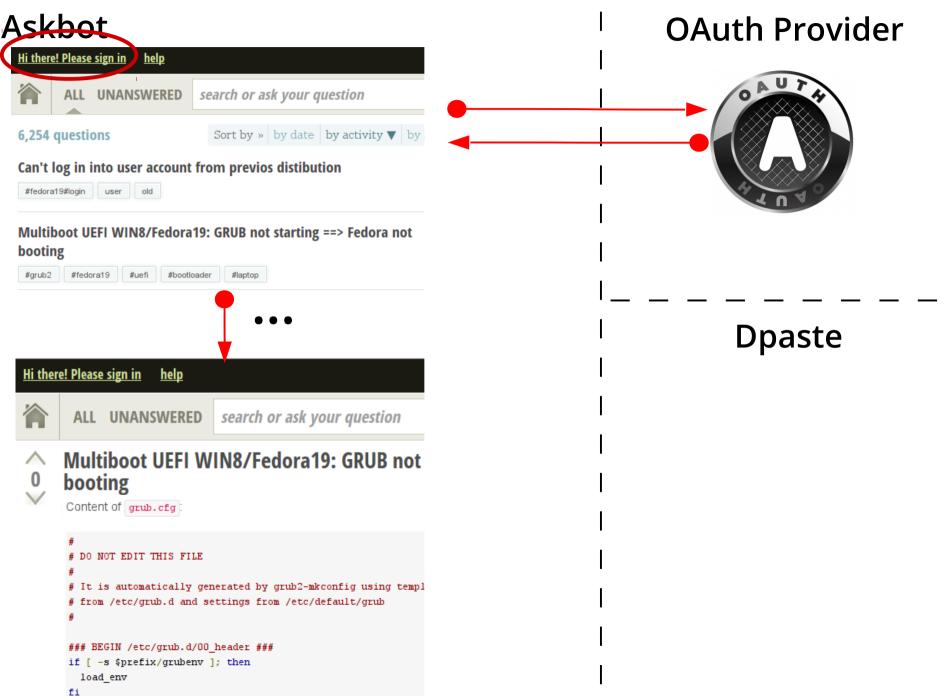
#grub2 #fedora19 #uefi #bootloader #laptop

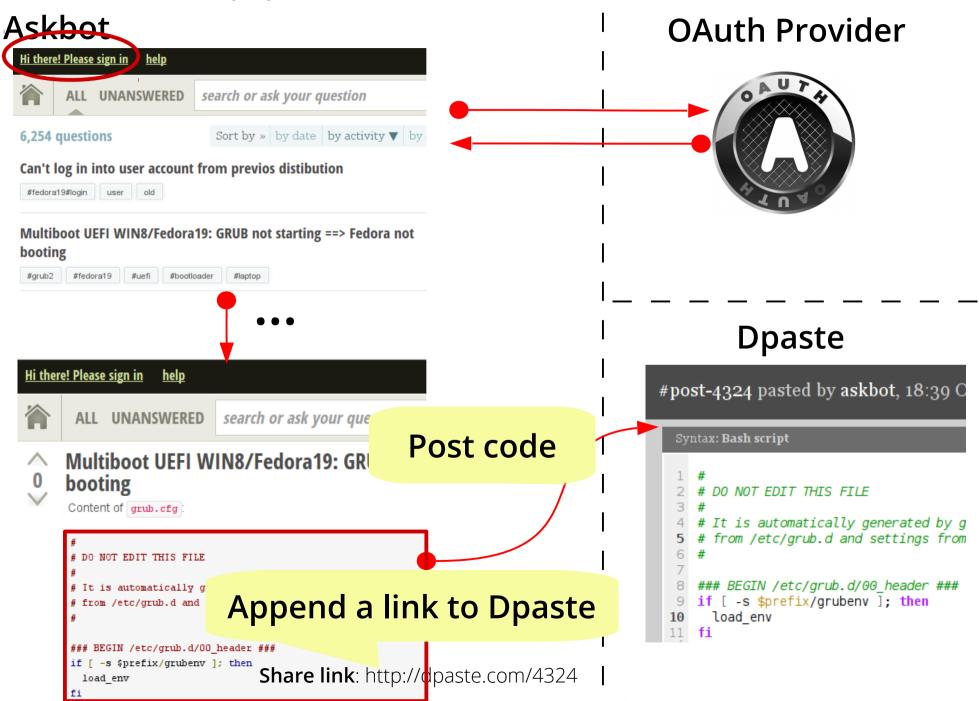
OAuth Provider

Dpaste

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6,254 qu				ate by activity ▼ by	
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Multiboo booting	ot UEFI WIN8/	Fedora19: 0	GRUB not starti	ing ==> Fedora not	
#grub2	#fedora19 #uefi	#bootloader	#laptop		

Dpaste







Askbot + OAuth + Dpaste = 183K LoC!

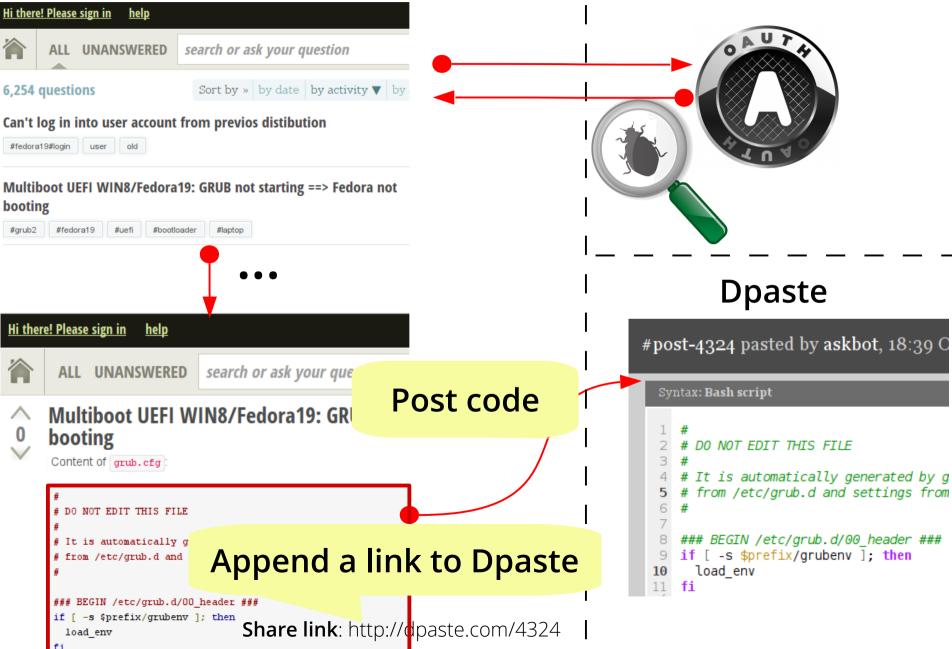
Aire can support large Django web applications

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				4	# It is automatically generated by g
					# from /etc/grub.d and settings from
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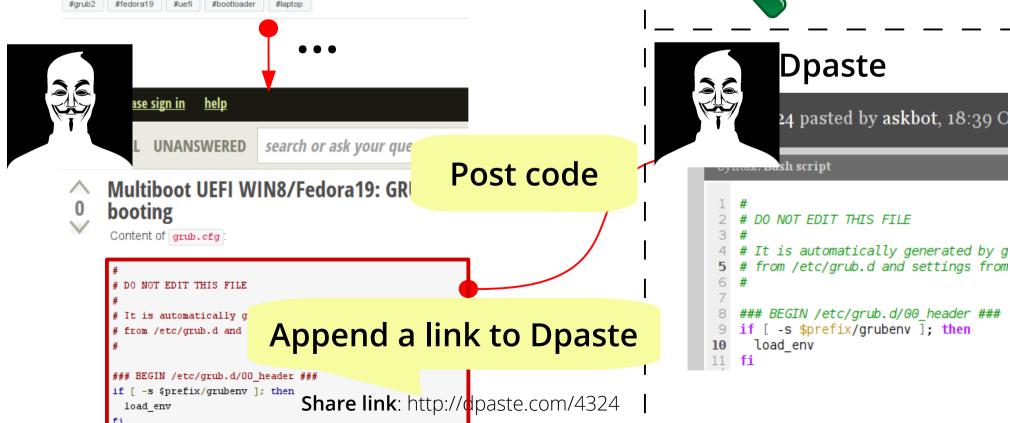
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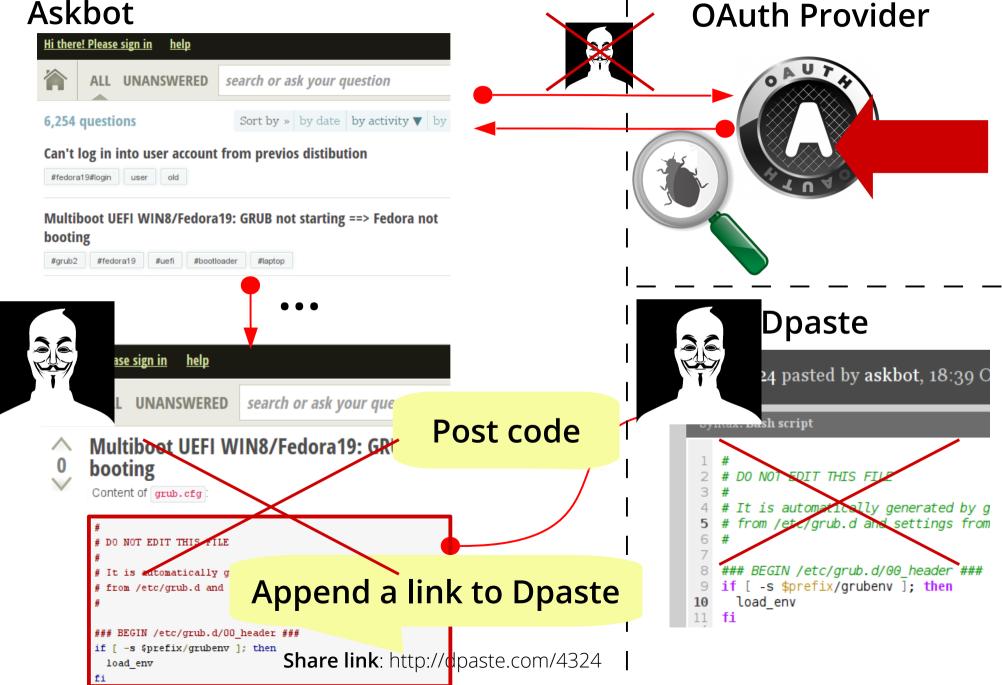


Aire enables automatic recovery Askbot **OAuth Provider** <u>Hi there! Please sign in</u> <u>ZE</u> help ALL UNANSWERED search or ask your question 6,254 questions Sort by » by date by activity ▼ by Can't log in into user account from previos distibution #fedora19#login user old Multiboot UEFI WIN8/Fedora19: GRUB not starting ==> Fedora not booting #fedora19 #bootloader #arub2 #uefi #laptop



Aire enables automatic recovery

Askbot



Aire enables automatic recovery

- Askbot, OAuth, and Dpaste are correctly recovered
 - Even when Dpaste is temporary unavailable
 - Even when Dpaste goes offline
- More examples in paper:
 - Intrusion recovery (synthetic)
 - Mistakes on ACL setting
 - Misconfigured versioning spreadsheet

Aire has moderate runtime overheads

Workload	Req/s without Aire	Req/s with Aire	Logs / req With Aire			
Reading	21.58	17.58	5.52 KB			
Writing	23.26	16.20	9.24 KB			
 19-30% throughput reduction 5-9KB/req storage overheads 						

→ Moderate overheads for websites which care integrity more than performance

Aire's repair is efficient

	Askbot	OAuth	DPaste
Repaired Reqs	105 / 2196	2/9	1 / 496
Remote repair reqs	1	1	0
Local repair time	84.06 sec	0.10 sec	3.91 sec
Normal exec. time	177.58 sec	0.01 sec	0.02 sec

- Experiment setting:
 - Attacker logins as a victim user and writes a post
 - 100 legitimate users post 5 questions and navigate
 - All users are affected by the attack (read attacker's post)

Aire's repair is Repair in Askbot propagates to OAuth and Dpaste

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Total repair takes **x2 shorter** than normal execution, although **x10 slower** in replaying a request for repair ost

- 100 legitimate users post 5 questions and navigate
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Related work

- Intrusion recovery with selective re-execution:
 - Retro [OSDI'10], Warp [SOSP'11]
 - → Use them as building blocks for **asynchronous** repair
- Intrusion recovery in distributed systems:
 - Heat-ray [SOSP'09], Polygraph [EuroSys'09], Dare [APsys'12]
 - → Automatic recovery in **loosely coupled web services**

Summary

- Aire recovers integrity of distributed web services
 - Define a **repair protocol**
 - Support asynchronous and decentralized repair
 - Propose partial repair consistency